

Code Injection (Process Injection)

Code Injection

- The technique of injecting malicious code into a target process's memory and executing the malicious code within the context of the target process is called code injection (or process injection).
- An attacker typically chooses a legitimate process (such as **explorer.exe** or **svchost.exe**) as the target process. Once the malicious code is injected into the target process, it can then perform malicious actions
- After injecting the code into the memory of the target process, the malware component responsible for injecting code can either continue to persist on the system, thereby injecting code into the target process every time the system reboots, or it can delete itself from the filesystem, keeping the malicious code in-memory only.

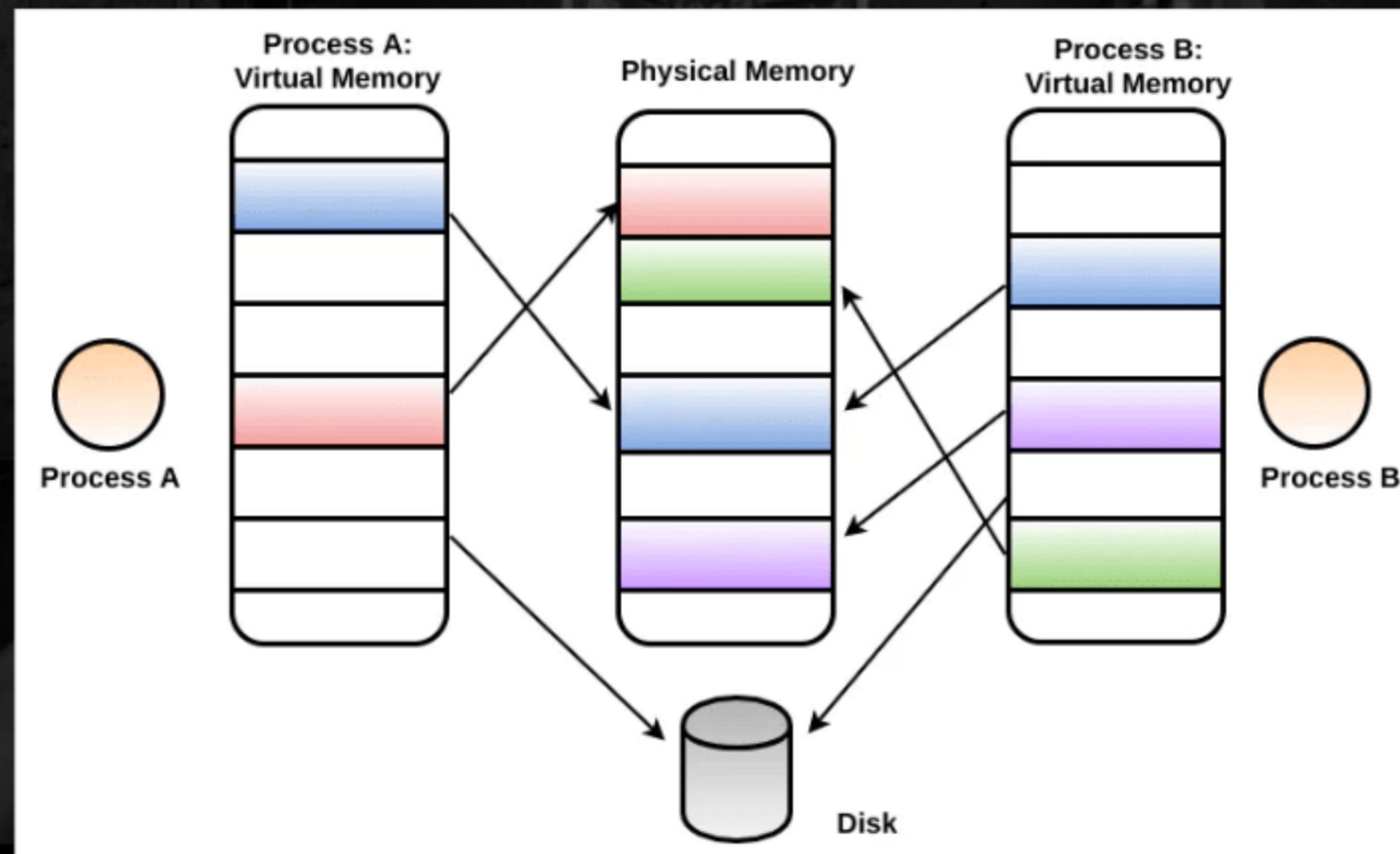
The injected code could be a module such as an executable, DLL, or even shellcode.

Once the code is injected into the remote process, an adversary can do the following things:

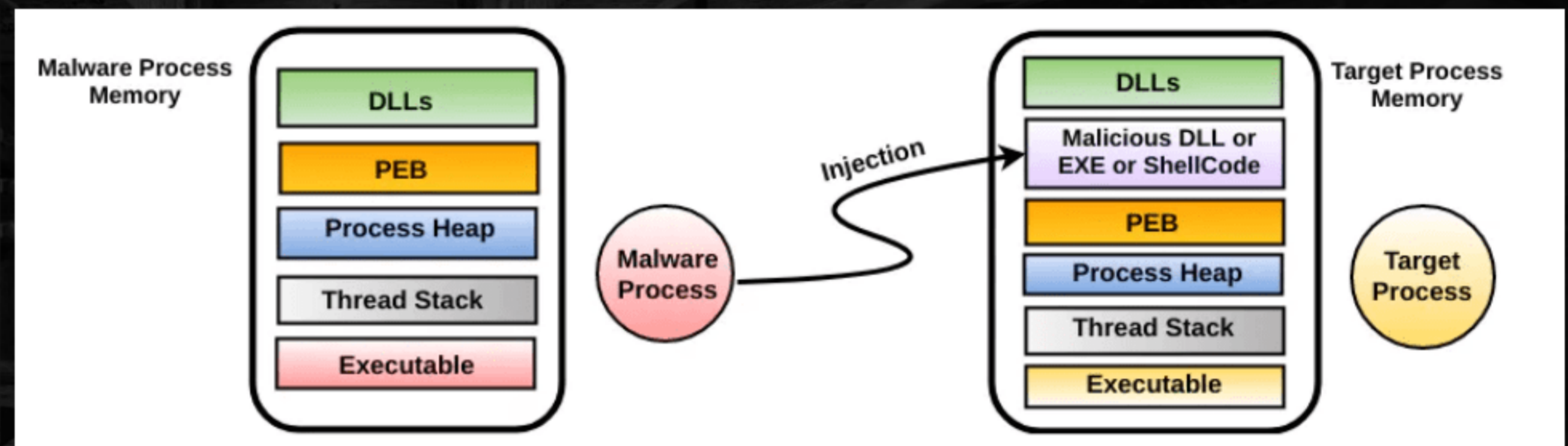
- Force the remote process to execute the injected code to perform malicious actions (such as downloading additional files or stealing keystrokes).
- Inject a malicious module (such as a DLL) and redirect the API call made by the remote process to a malicious function in the injected module. The malicious function can then intercept the input parameters of the API call, and also filter the output parameters.
- Injecting code into an already running process allows an adversary to achieve persistence.
- Injecting code into trusted processes allows an attacker to bypass security products (such as whitelisting software) and hide from the user.

Virtual Memory

Windows operating system provides each new process that is created, with its own private memory address space (called the **process memory**). The process memory is a part of virtual memory; virtual memory is not real memory, but an illusion created by the operating system's memory manager.



The following screenshot should give you a high-level overview of code injection techniques in the user-space:





Code Injection Techniques

Remote Executable/Shellcode Injection

- The malicious code is injected into the target process memory directly, without dropping the component on the disk.
- The injected malicious code is forced to execute by creating a remote thread via **CreateRemoteThread()**, and the start of the thread is made to point to the code/function within the injected block of code.
- The advantage of this method is that the malware process does not have to drop the malicious DLL on the disk; it can extract the code to inject from the resource section of the binary, or get it over the network and perform code injection directly.

Demo 11 - Remote Executable Injection

Remote DLL Injection

Remote DLL Injection

In this injection technique, the malware process creates a thread in the target process and the thread is made to call **LoadLibrary()** by passing a malicious DLL path as the argument. Since the thread gets created in the target process, the target process loads the malicious DLL into its address space. Once the target process loads the malicious DLL, the operating system automatically calls the DLL's **DllMain()** function, thus executing the malicious code.

Step 1:

The malware process identifies the target process and opens a handle to the target process. In the following screenshot, the malware calls **OpenProcess()** by passing the pid of **explorer.exe** (**0x624**, which is **1572**) as the third parameter. The return value of **OpenProcess()** is the handle to the **explorer.exe** process:

0040143D	53	push ebx	EAX 0000003C '<'
0040143E	6A 00	push 0	EBX 00000624 L' A'
00401440	68 FF FF 1F 00	push 1FFFFFFF	Default (stdcall)
EIP → 00401445	FF 15 10 A0 40 00	call dword ptr ds:[&OpenProcess]	1: [esp] 001FFFFFFF
0040144B	8B F8	mov edi, eax	2: [esp+4] 00000000
0040144D	85 FF	test edi, edi	3: [esp+8] 00000624 ←
0040144F	75 1E	jne nps.40146F	

Step 2:

Malware process allocates memory in the target process using the *VirtualAllocEx()* API. The **1st** argument (**0x30**) is the handle to *explorer.exe* (the target process). The **3rd** argument, **0x27 (39)**, represents the number of bytes to be allocated in the target process, and the **5th** argument (**0x4**) is a constant value that represents the memory protection of *PAGE_READWRITE*. The return value of *VirtualAllocEx()* is the address of the allocated memory in *explorer.exe*

0040148C	6A 04	push 4	
0040148E	68 00 10 00 00	push 1000	
00401493	56	push esi	
00401494	6A 00	push 0	
00401496	57	push edi	
EIP → 00401497	FF 15 58 A0 40 00	call dword ptr ds:[&VirtualAllocEx]	
0040149D	8B D8	mov ebx, eax	
0040149F	85 DB	test ebx, ebx	
004014A1	75 1E	jne nps.4014C1	

EAX	0000005F	' '
EBX	00000624	L'â'
Default (stdcall)		
1:	[esp]	00000030 ←
2:	[esp+4]	00000000
3:	[esp+8]	00000027 ←
4:	[esp+C]	00001000
5:	[esp+10]	00000004 ←

Step 3:

Malware uses `WriteProcessMemory()` to copy the DLL pathname to the allocated memory in the target process. The *2nd* argument, `0x01E30000`, is the address of the allocated memory in the target process, and the *3rd* argument is the full path to the DLL that will be written to the allocated memory address `0x01E30000` in `explorer.exe`

56	<code>push esi</code>	EAX	0012F674
FF B5 DC FB FF FF	<code>push dword ptr ss:[ebp-424]</code>	EBX	01E30000
53	<code>push ebx</code>	Default (stdcall)	
57	<code>push edi</code>	1: [esp]	00000030
FF 15 4C A0 40 00	<code>call dword ptr ds:[<&WriteProcessMemory>]</code>	2: [esp+4]	01E30000
85 C0	<code>test eax, eax</code>	3: [esp+8]	0012FABC "C:\\Users\\test\\AppData\\Roaming\\adpr.dll"
75 2A	<code>jne nps.401505</code>	4: [esp+C]	00000027

Step 4:

Malware determines the address of **LoadLibrary()** in **kernel32.dll**; to do that, it calls the **GetModuleHandleA()** API and passes **kernel32.dll** as the argument, which will return the base address of **Kernel32.dll**. Once it gets the base address of **kernel32.dll**, it determines the address of **LoadLibrary()** by calling **GetProcAddress()**

Step 5:

Malware calls **CreateRemoteThread()** which creates a thread in the target process. The **1st** argument, **0x30**, to **CreateRemoteThread()** is the handle to the **explorer.exe** process, in which the thread will be created. The **4th** argument is the address in the target process memory where the thread will start executing, which is the address of **LoadLibrary()**, and the **5th** argument is the address in the target process memory that contains the full path to the DLL.

004015D9	50	push eax	EAX 00000000
004015DA	50	push eax	EBX 01E30000
004015DB	53	push ebx	Default (stdcall)
004015DC	56	push esi	1: [esp] 00000030 ←
004015DD	50	push eax	2: [esp+4] 00000000
004015DE	50	push eax	3: [esp+8] 00000000
004015DF	57	push edi	4: [esp+C] 7791DE15 <kernel32.LoadLibraryA> ←
EIP → 004015E0	FF 15 30 A0 40 00	call dword ptr ds:[&CreateRemoteThread]	5: [esp+10] 01E30000 ←
004015E6	8B F0	mov esi, eax	6: [esp+14] 00000000 ←



DLL Injection Using APC (APC Injection)

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- The APC injection technique is similar to remote DLL injection, but instead of using **CreateRemoteThread()**, a malware makes use of **Asynchronous Procedure Calls (APCs)** to force the thread of a target process to load the malicious DLL.
- An APC is a function that executes asynchronously in the context of a particular thread. Each thread contains a queue of APCs that will be executed when the target thread enters an alertable state.
- A thread enters an alertable state if it calls one of the following functions:

```
SleepEx(),  
SignalObjectAndWait()  
MsgWaitForMultipleObjectsEx()  
WaitForMultipleObjectsEx()  
WaitForSingleObjectEx()
```

- The way APC injection technique works is, a malware process identifies the thread in the target process (the process in which the code will be injected) that is in an alertable state, or likely to go into an alertable state.
- It then places the custom code in that thread's APC queue by using the **QueueUserAPC()** function.
- The idea of queueing the custom code is that, when the thread enters the alertable state, the custom code gets picked up from the APC queue, and it gets executed by the thread of the target process.

Steps In APC Injection

This technique starts with the same four steps as **Remote DLL injection**:

Step 1: It opens a handle to the target process

Step 2: Allocates memory in the target process

Step 3: Copies the malicious DLL pathname into the allocated memory

Step 4: determines the address of **Loadlibrary()**

Step 5:

It opens a handle to the thread of the target process using the **OpenThread()** API. In the following example, The **3rd** argument, **0xBEC(3052)**, is the thread ID (TID) of the **iexplore.exe** process. The return value of **OpenThread()** is the handle to the thread of **iexplore.exe**:

0128827E	57	push edi	EAX	00000001
0128827F	6A 00	push 0	EBX	00000000
01288281	68 FF 03 1F 00	push 1F03FF	Default (stdcall)	
EIP → 01288286	FF 15 00 D4 4A 01	call dword ptr ds:[<&OpenThread>]	1: [esp]	001F03FF
0128828C	A3 68 09 55 01	mov dword ptr ds:[1550968],eax	2: [esp+4]	00000000
01288291	85 C0	test eax,eax	3: [esp+8]	00000BEC ← Thread Id of iexplore.exe

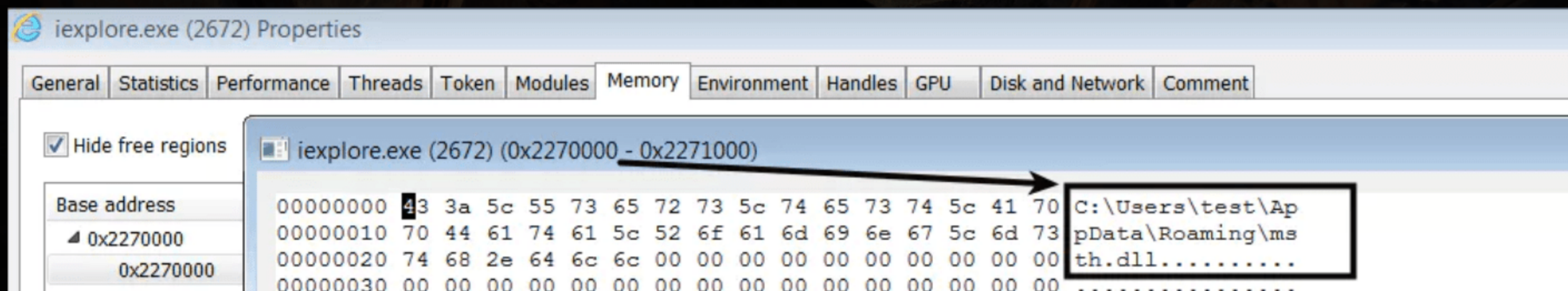
Step 6:

The malware calls **QueueUserAPC()** to queue the APC function. In the following screenshot, the **1st argument** is the pointer to the APC function that the malware wants the target thread to execute. In this case, the APC function is the **LoadLibrary()** whose address was determined previously. The **2nd argument**, **0x22c**, is the handle to the target thread of **iexplore.exe**. The **3rd argument**, **0x2270000**, is the address in the target process (**iexplore.exe**) memory containing the full path to the malicious DLL

01288A5F	FF 75 FC	push dword ptr ss:[ebp-4]
01288A62	FF 75 E8	push dword ptr ss:[ebp-18]
01288A65	FF 75 F4	push dword ptr ss:[ebp-C]
EIP → 01288A68	FF 15 10 D4 4A 01	call dword ptr ds:[<&QueueUserAPC>]
01288A6E	85 C0	test eax, eax
01288A70	75 2C	jne rmdi.1288A9E
01288A72	FF 15 88 D3 4A 01	call dword ptr ds:[<&GetLastError>]

EAX	00000015
EBX	00000000
Default (stdcall)	
1: [esp]	7791DE15 <kernel32.LoadLibraryA>
2: [esp+4]	0000022C
3: [esp+8]	02270000

The following screenshot shows the content of the address 0x2270000 in Internet Explorer's process memory (this was passed as the **3rd argument** to **QueueUserAPC()**); this address contains the full path to the DLL that was previously written by the malware:



At this point the injection is complete and when the thread of the target process enters alertable state, the thread executes **LoadLibrary()** from the **APC queue** by passing the full path to the DLL as the argument to **LoadLibrary()**, as a result the malicious DLL gets loaded into target process address space which in turn invokes the **DLLMain()** function containing the malicious code

DLL Injection Using SetWindowsHookEx()

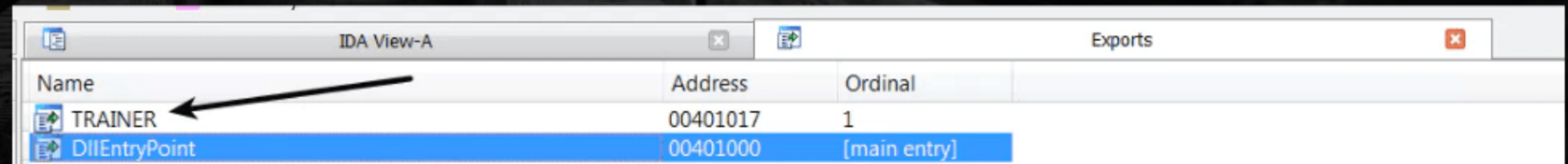
DLL Injection Using SetWindowsHookEx()

- The **SetWindowsHookEx()** API can also be used to load a DLL into a target process address space and execute malicious code.
- A malware first loads the malicious DLL into its own address space. It then installs a hook procedure (a function exported by the malicious DLL) for a particular event (such as a keyboard or mouse event), and it associates the event with the thread of the target process (or all of the threads in the current desktop).
- When a particular event is triggered, for which the hook is installed, the thread of the target process will invoke the hook procedure. To invoke a hook procedure defined in the DLL, it must load the DLL (containing the hook procedure) into the address space of the target process.

Example

The following describes the steps performed by the malware sample (**Trojan Padador**) :

Step 1: The malware executable drops a DLL named **tckdll.dll** on the disk. The DLL contains an **entrypoint** function, and an export function named **TRAINER**. The DLL entry point function does not do much, whereas the **TRAINER** function contains the malicious code.



Name	Address	Ordinal
TRAINER	00401017	1
DllEntryPoint	00401000	[main entry]

Step 2:

Malware loads the DLL (*tckdll.dll*) into its own address space using the *LoadLibrary()* API, but no malicious code is executed at this point. The return value of *LoadLibrary()* is the handle to the loaded module (*tckdll.dll*). It then determines the address of the *TRAINER* function by using *GetProcAddress()*:

00401047	push	offset LibFileName ; "tckdll.dll"	← Loads tckdll.dll into its own address space
0040104C	call	LoadLibraryA	
00401051	mov	hmod, eax	← Determines the address of TRAINER function
00401056	push	offset ProcName ; "TRAINER"	
0040105B	push	eax ; hModule	
0040105C	call	GetProcAddress	

Step 3:

Malware uses the handle to tckdll.dll and the address of the TRAINER function to register a hook procedure for the keyboard event. The 1st argument, WH_KEYBOARD (constant value 2), specifies the type of event that will invoke the hook routine. The 2nd argument is the address of the hook routine, which is the address of the TRAINER function. The 3rd argument is the handle to the tckdll.dll, which contains the hook procedure. The fourth argument, 0, specifies that the hook procedure must be associated with all of the threads in the current desktop.

```
0040105C  call  GetProcAddress
00401061  push  0           ; dwThreadId
00401063  push  hmod       ; hmod
00401069  push  eax        ; lpfn
0040106A  push  WH_KEYBOARD ; idHook
0040106C  call  SetWindowsHookExA
```

← Registers hook procedure for the keyboard event

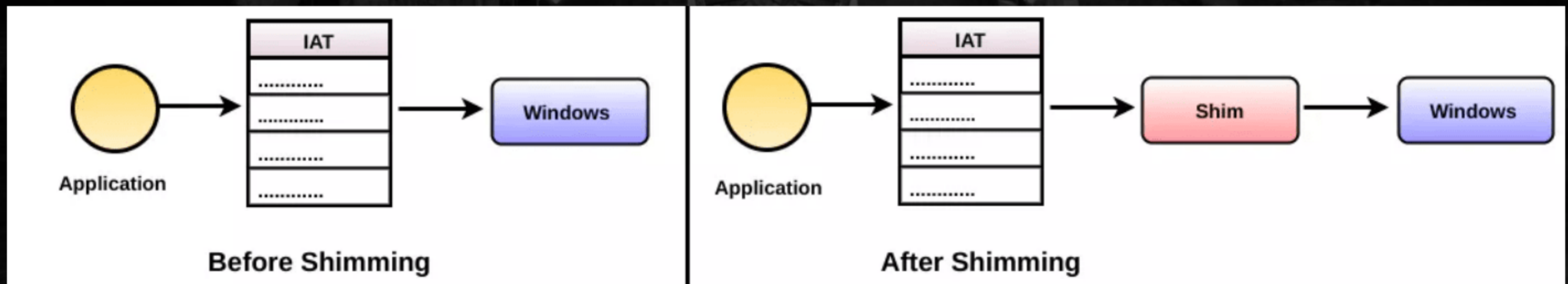
After these steps, when the keyboard event is triggered within an application, that application will load the malicious DLL and invoke the TRAINER function.

DLL Injection Using The Application Compatibility Shim

DLL Injection Using The Application Compatibility Shim

- The Microsoft Windows application compatibility infrastructure/framework (application shim) is a feature that allows programs created for older versions of the operating system (such as Windows XP) to work with modern versions of the operating system (such as Windows 7 or Windows 10).
- The shims are provided by Microsoft to the developers so that they can apply fixes to their programs without rewriting the code. When a shim is applied to a program, and when the shimmed program is executed, the shim engine redirects the API call made by the shimmed program to shim code; this is done by replacing the pointer in the IAT with the address of the shim code.

The following screenshot should help you to understand the differences in interactions between the normal and shimmed applications in the Windows operating system:



Suppose that a few years back (before the release of Windows 7), you wrote an application (**xyz.exe**) which checked the OS version, before performing some useful operation. Let's suppose that your application determined the OS version by calling the **GetVersion()** API in **kernel32.dll**. In short, the application did something useful only if the OS version was Windows XP. Now, if you take that application (**xyz.exe**) and run it on Windows 7, it will not do anything useful, because the OS version returned on Windows 7 by **GetVersion()** does not match with Windows XP. To make that application work on Windows 7, you can either fix the code and rebuild the program, or you can apply a shim called **WinXPVersionLie** to that application (**xyz.exe**).

Demo 12 - Injecting DLL Using Shim



Hollow Process Injection (Process Hollowing)

Hollow Process Injection (Process Hollowing)

- Process hollowing, or Hollow Process Injection, is a code injection technique in which the executable section of the legitimate process in the memory, is replaced with a malicious executable.
- This technique allows an attacker to disguise his malware as a legitimate process and execute malicious code.
- The advantage of this technique is that the path of the process being hollowed out will still point to the legitimate path, and, by executing within the context of a legitimate process, the malware can bypass firewalls and host intrusion prevention systems.



Demo 13 - Hollow Process Injection