

Apply Cryptographic Algorithms



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Cryptography for SSCP®

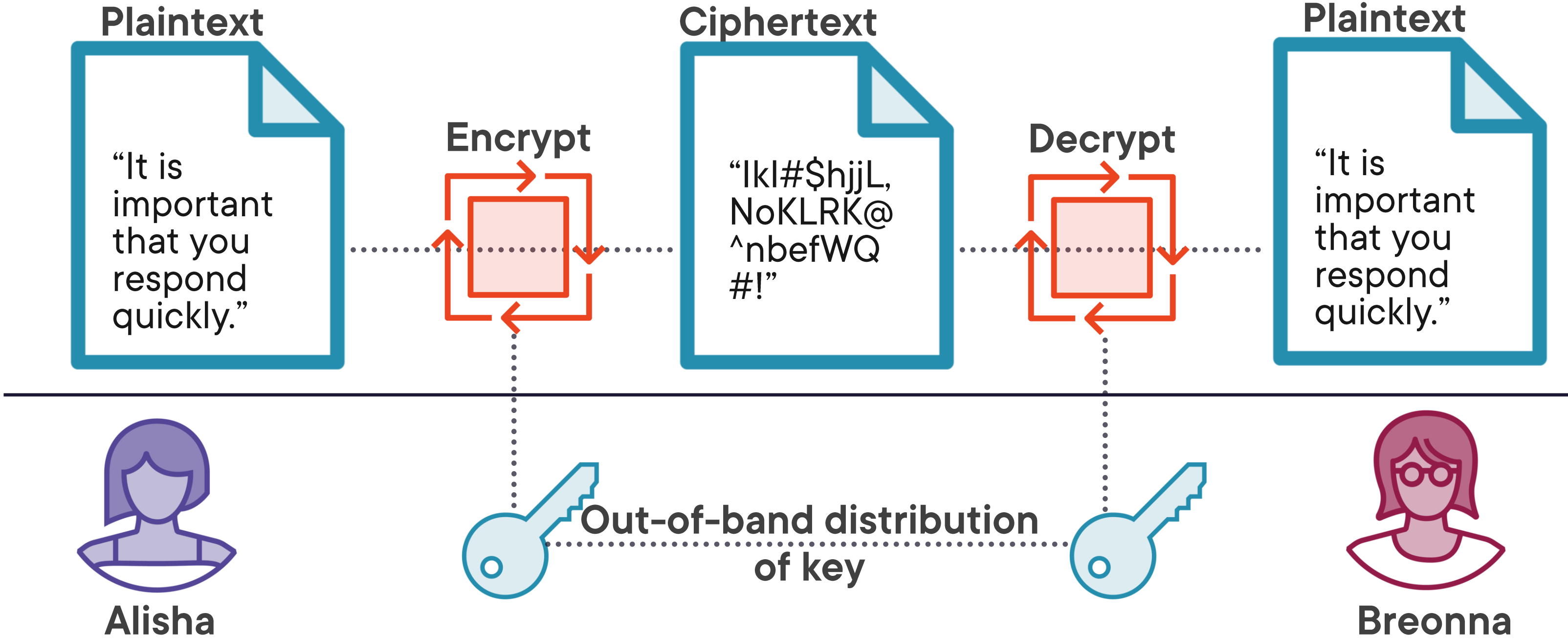
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Symmetric Algorithms Overview



Symmetric Key Process Flow



Foundational Math for Symmetric Ciphers

1 ⊕ 1 = 0

0 ⊕ 0 = 0

1 ⊕ 0 = 1

0 ⊕ 1 = 1



Symmetric Algorithms Issues

Advantages

Excellent for confidentiality

Encryption and decryption are relatively fast

Generally patent free use without cost

Provides some level of integrity (HMAC, Keyed-hash)

Disadvantages

Key distribution challenges

Poor authentication and no non-repudiation



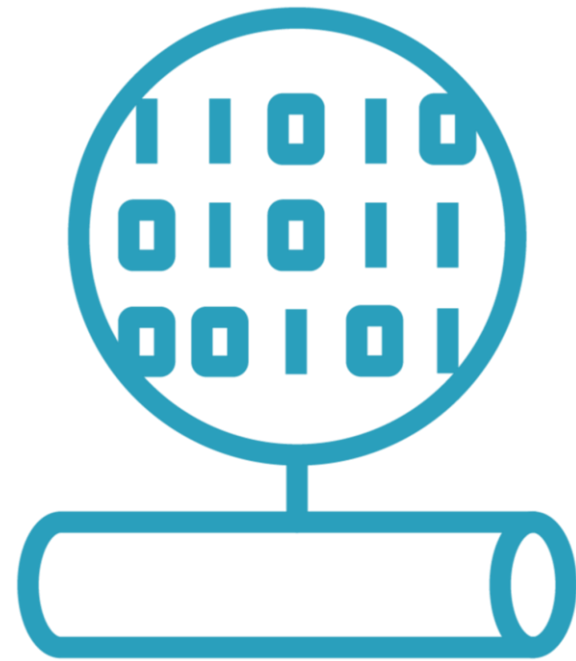
**The problem with
symmetric keys:
 $n(n-1)/2$**



5,000 users = 12,497,500 keys

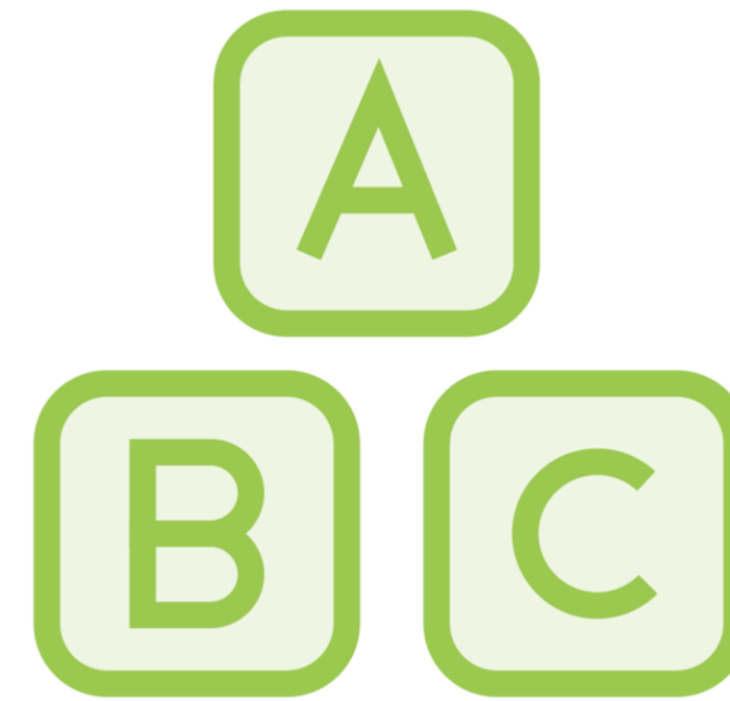


Symmetric Cipher Types



Stream-based

Encryption operations is on a constant stream of 0s and 1s.



Block-based

Encryption operation is on fixed blocks of plaintext.



Symmetric Stream-based Algorithm Types and Characteristics



Symmetric Stream-based Algorithms Applicability and Identification

Data in transit

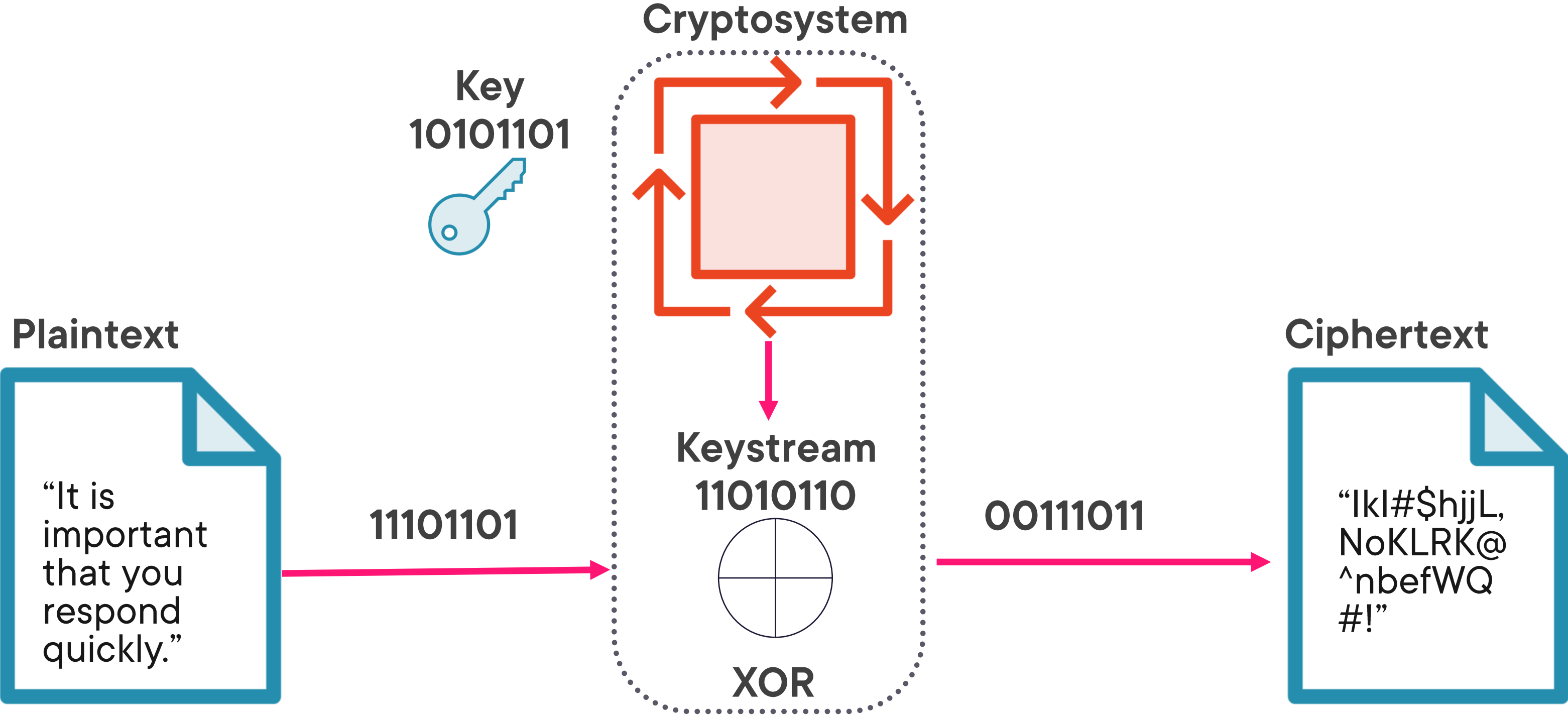
High-speed with minimal latency

Often embedded in hardware

Each bit or byte is encrypted



Symmetric Stream-based Basics



RC4
SEAL
WAKE
A5

Symmetric Stream-based Algorithms



RC4 Characteristics

Published in 1987

Variable key length 1-256 bytes

Key initializes a state vector
which generates keystream to
XOR plaintext

Considered to be deprecated by
NIST



The Initial Block-based Symmetric Algorithm



Symmetric Block- based Algorithms Applicability and Identification

**Encrypts fixed plaintext data blocks of 64, 128,
192, 256 bits**

Viewed as more robust than stream ciphers

May have modes that behave as stream



Horst Feistel headed up research at IBM in the 1960's that eventual led to the release of Data Encryption Standard (DES) in 1977



64-bit blocks
Two 32-bit halves
56-bit key
48-bit subkeys
16 rounds of encryption

Data Encryption Standard (DES)



Five Modes of DES

**Electronic Code
Book (ECB)**

Block mode

**Cipher Block
Chaining (CBC)**

Block mode / IV

**Cipher Feedback
(CFB)**

Stream mode

**Output Feedback
(OFB)**

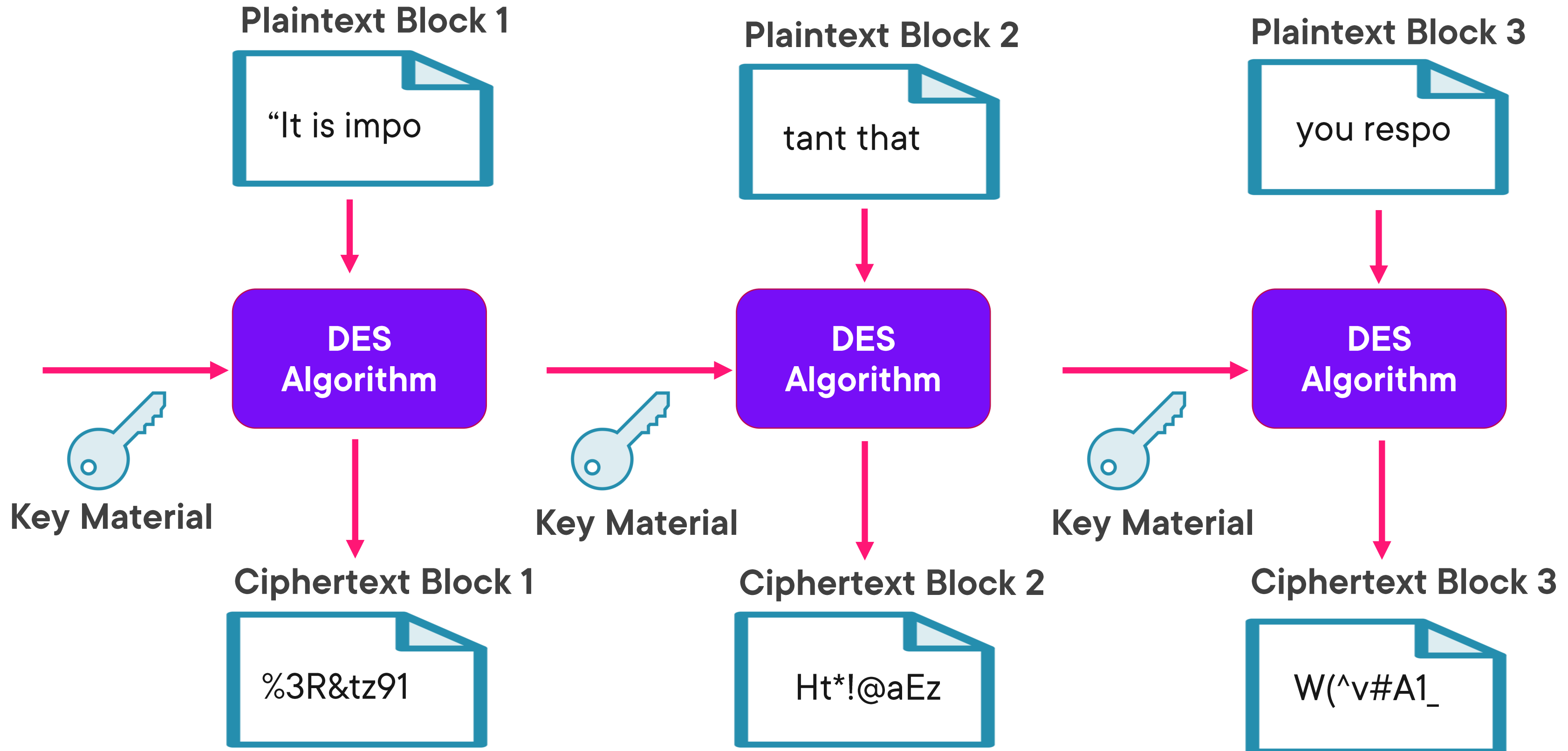
Stream mode

**Counter Mode
(CTR)**

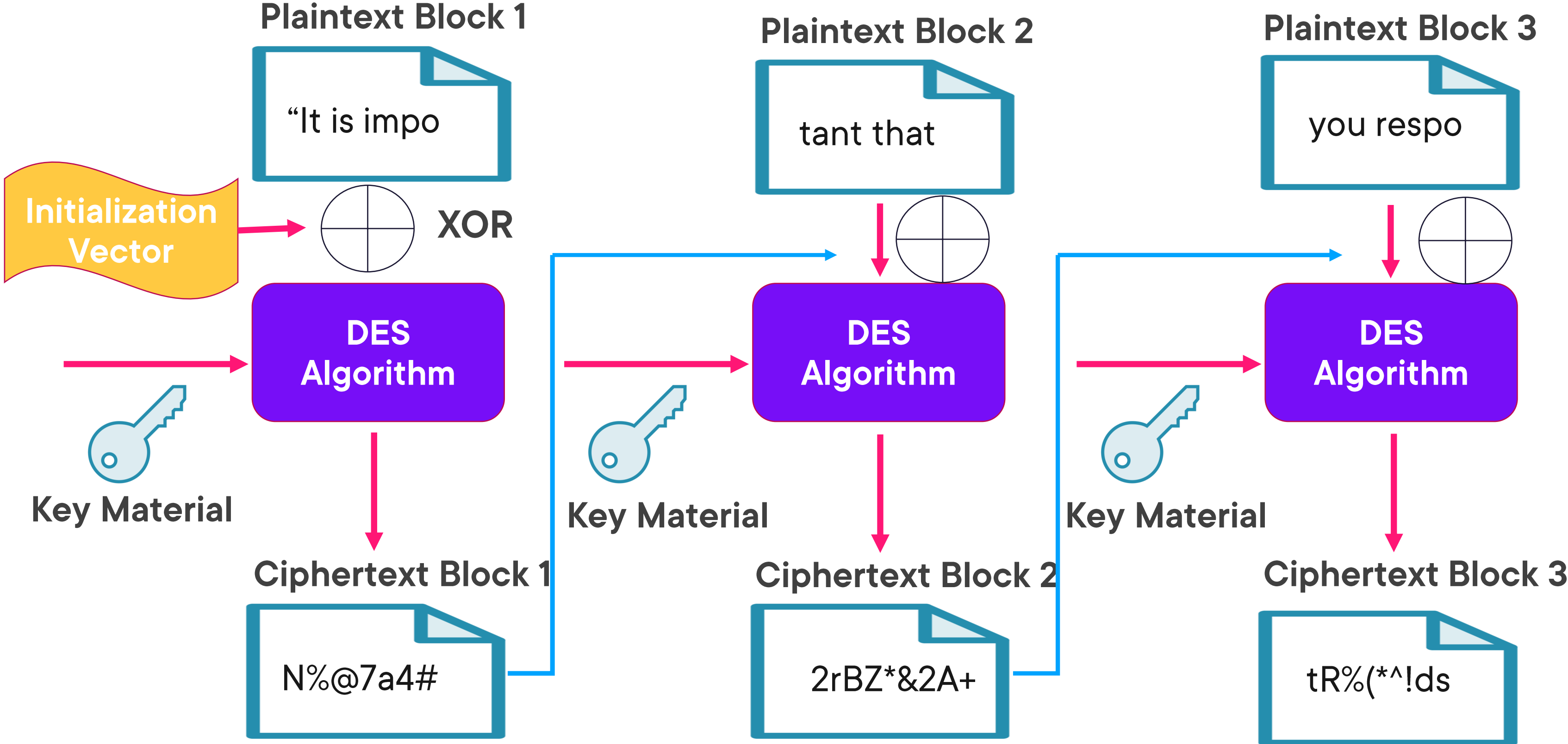
Stream mode



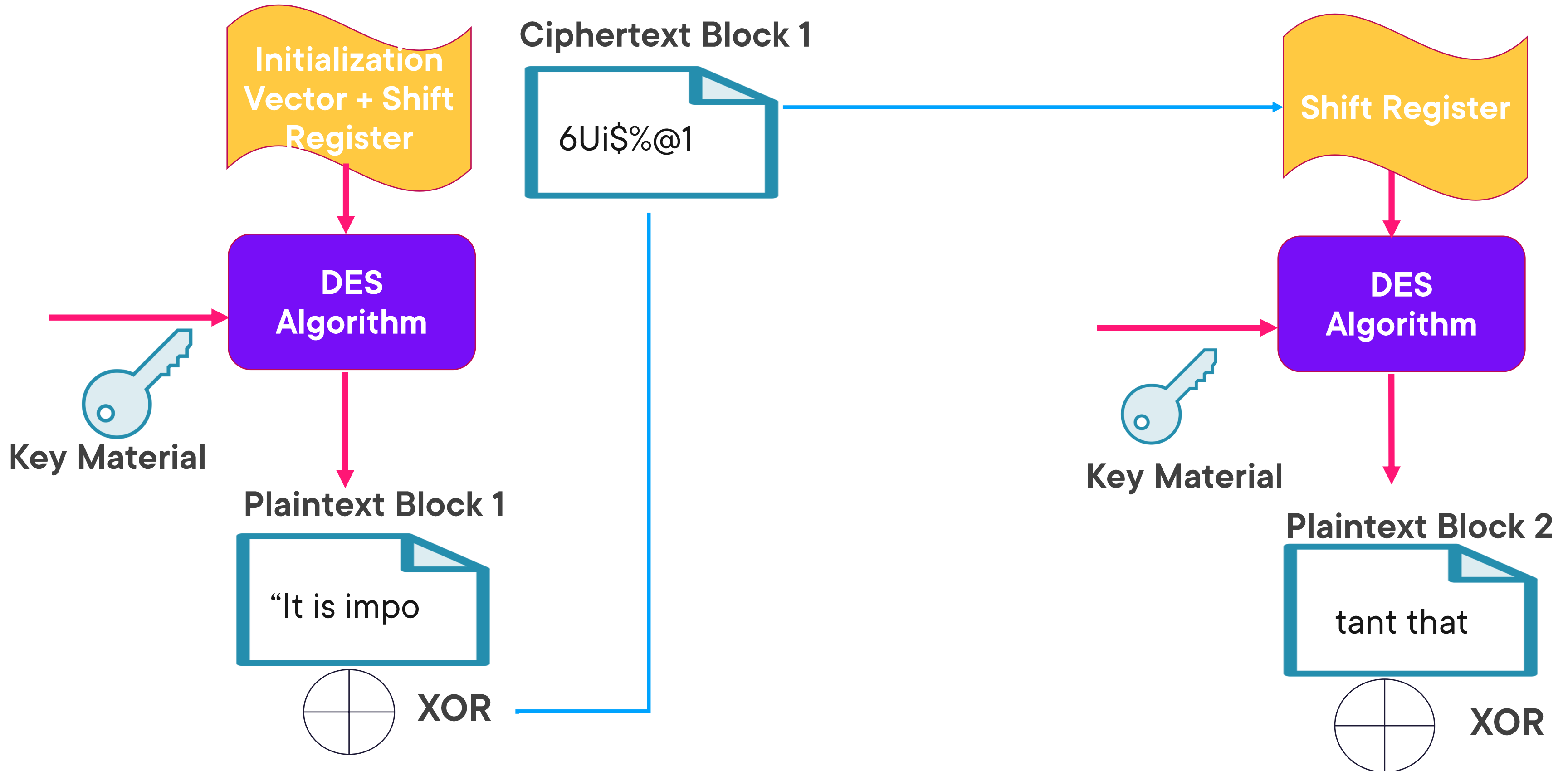
DES Block Mode Electronic Code Book (ECB)



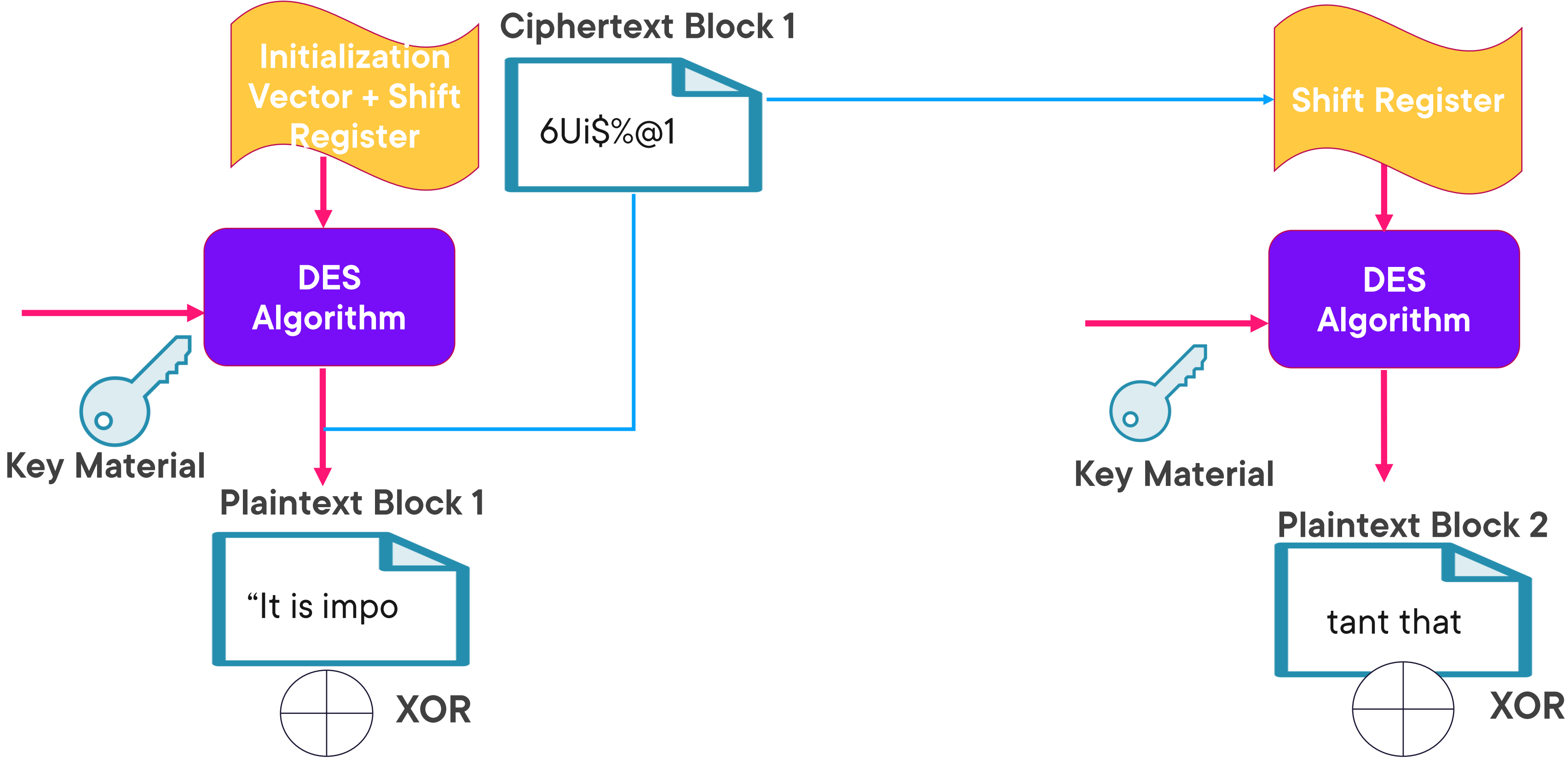
DES Block Mode Cipher Block Chaining (CBC)



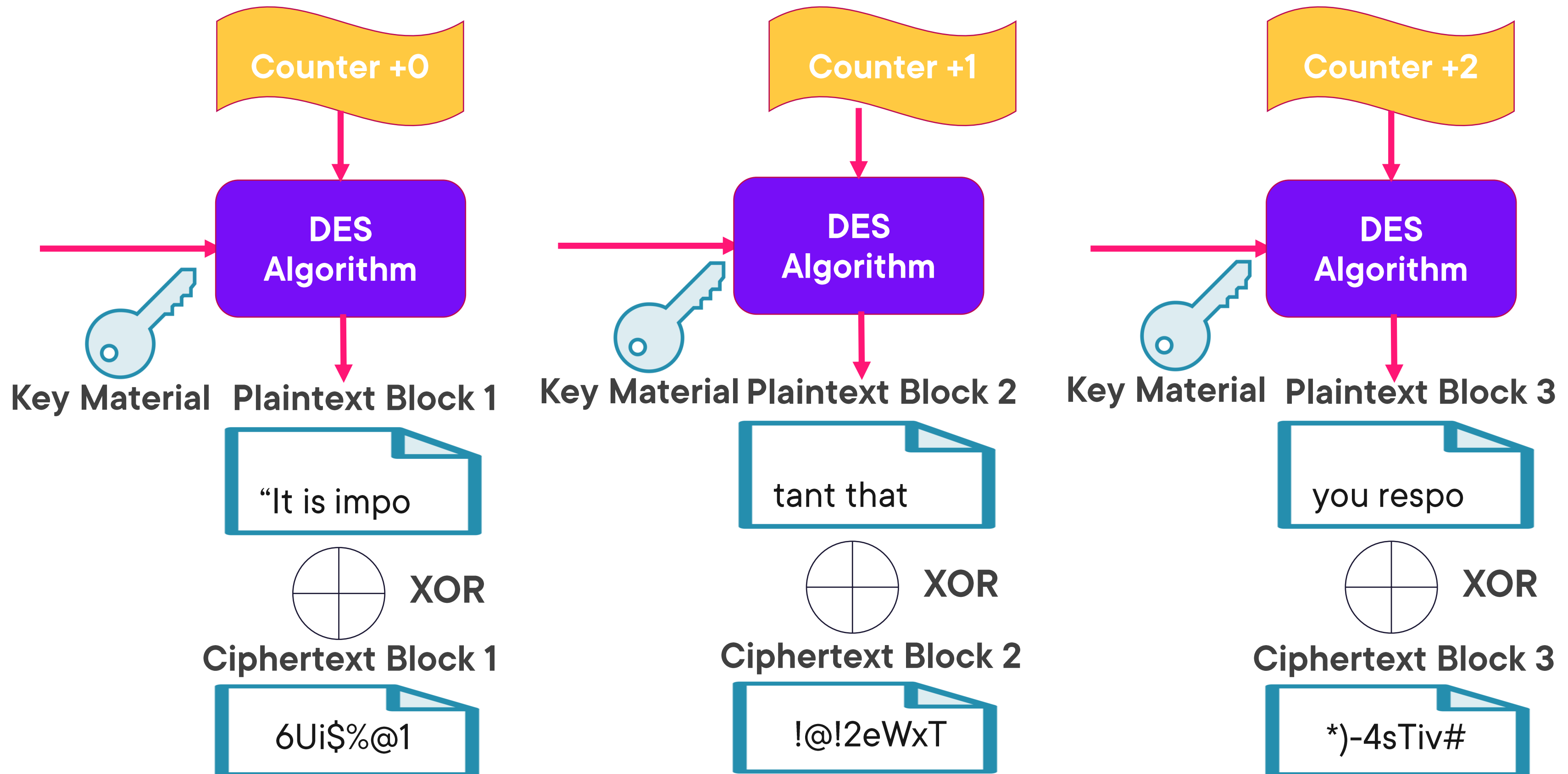
DES Stream Mode Cipher Feedback (CFB)



DES Stream Mode Output Feedback (OFB)



DES Stream Mode Counter (CTR)



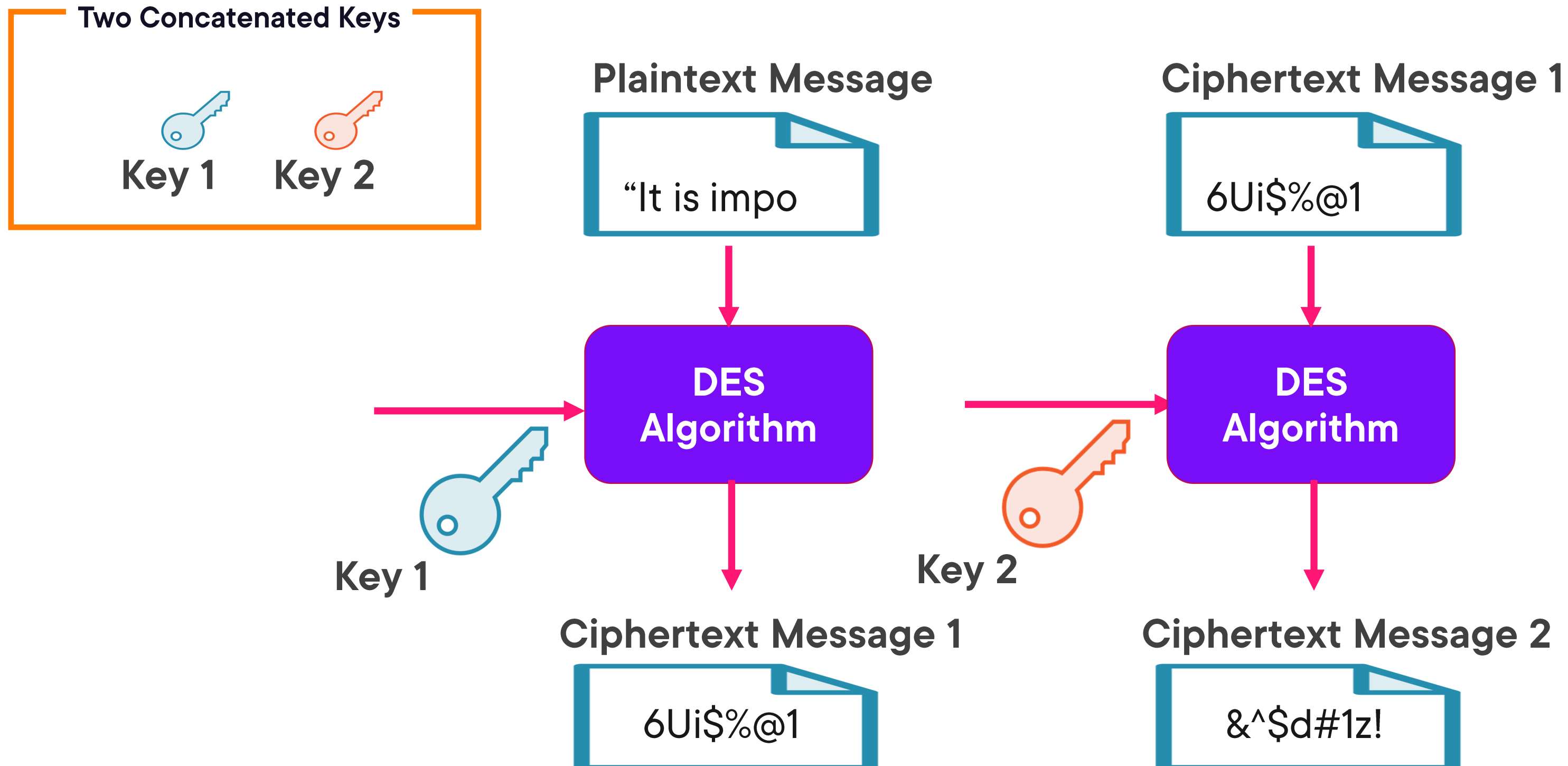
Double (2DES) and Triple DES (3DES)



For its time DES served the purpose of confidentiality, but Moore's Law and cryptanalytic advances exposed weaknesses.



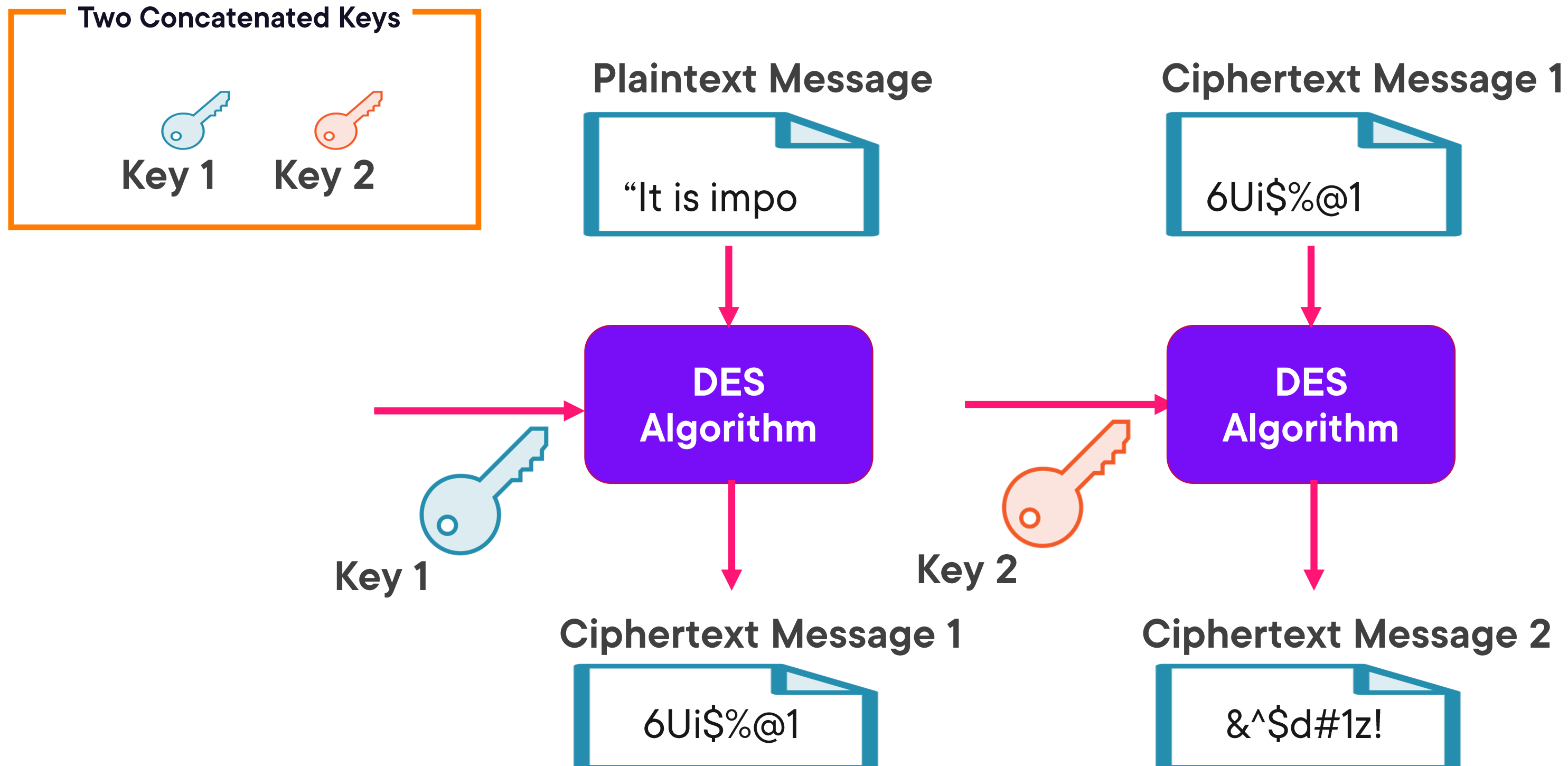
Double DES (2DES)



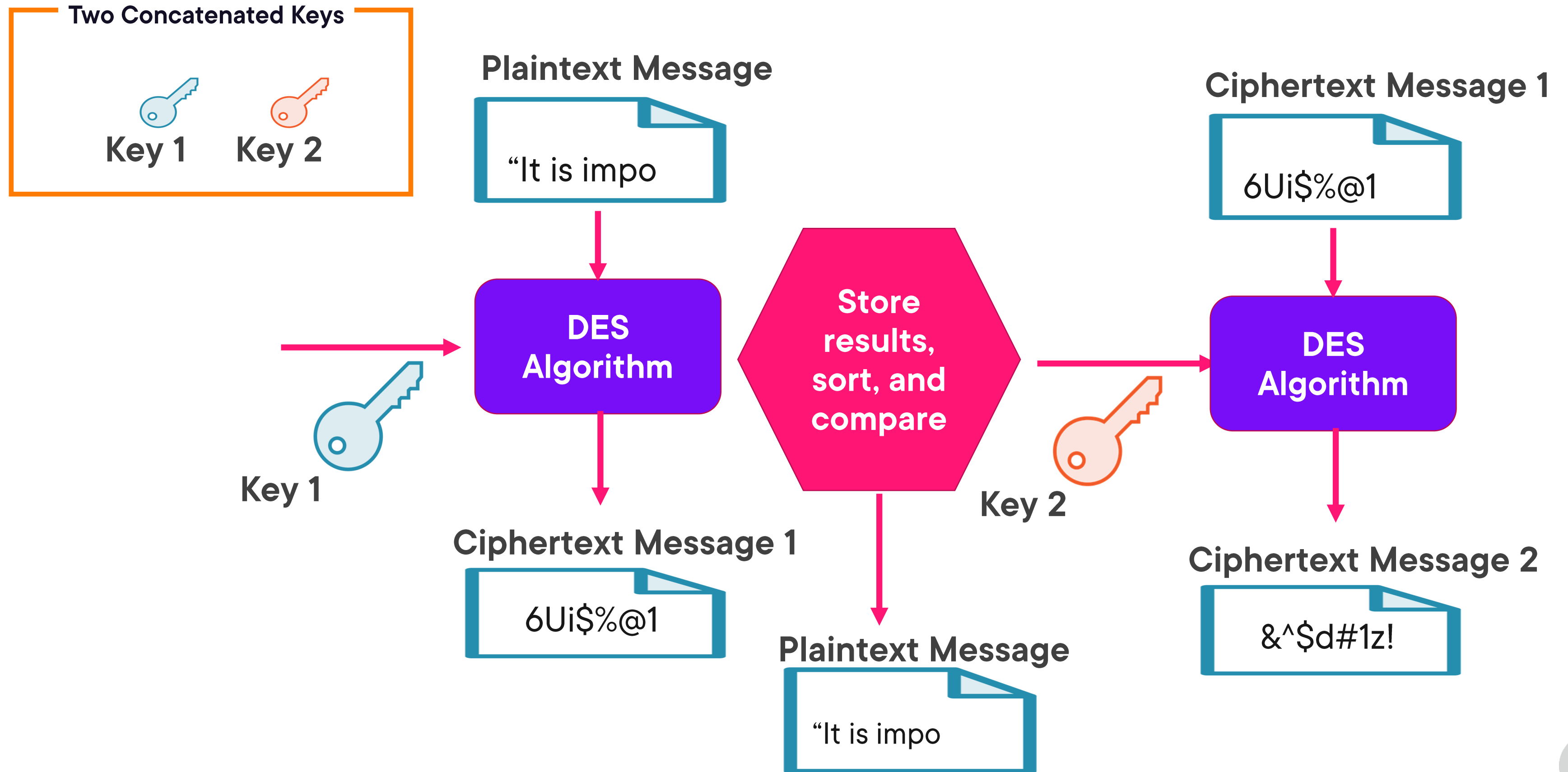
Double DES calculates $C = EK_2(EK_1(P))$



Meet-in-the-Middle Attack



Meet-in-the-Middle Attack



**Because of Meet-in-the-Middle
Double DES 2^{112} only provides 2^{57}
relative strength.**



Triple DES (3DES) EEE2 or EDE2

Two Concatenated
Keys



Key 1



Key 2

Plaintext Message

"It is impo

Ciphertext Message 1

6Ui\$%@1

Ciphertext Message 2

&^\$d#1z!

DES
Algorithm

DES
Algorithm

DES
Algorithm

Key 1

Key 2

Key 1

Ciphertext Message 1

Ciphertext Message 2

Ciphertext Message 3

6Ui\$%@1

&^\$d#1z!

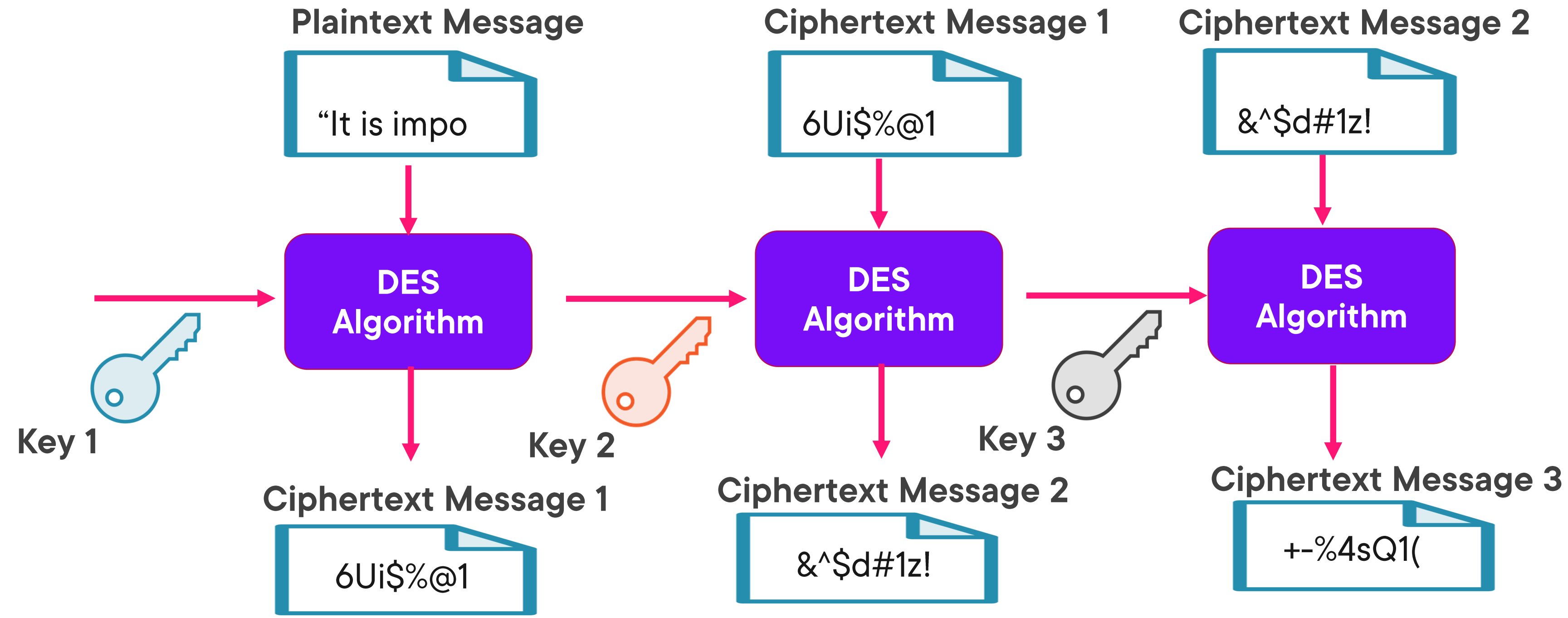
+-%4sQ1(



Three Concatenated Keys

Key 1 Key 2 Key 3

Triple DES (3DES) EEE3 or EDE3



Triple DES calculates

$C = EK_1(EK_2(EK_1(P)))$ for the EEE2 mode

$C = EK_1(DK_2(EK_1(P)))$ for the EDE2 mode

$C = EK_3(EK_2(EK_1(P)))$ for the EEE3 mode

$C = EK_3(DK_2(EK_1(P)))$ for the EDE3 mode



**Because of Meet-in-the-Middle
Triple DES 2^{168} only provides 2^{112} relative
strength.**



Symmetric Block-based Algorithm Types and Characteristics



“Beginning in 1997, NIST worked with industry and the cryptographic community to develop an Advanced Encryption Standard (AES). The overall goal was to develop a Federal Information Processing Standard (FIPS) specifying an encryption algorithm capable of protecting sensitive government information well into the 21st century. The algorithm was expected to be used by the U.S. Government and, on a voluntary basis, by the private sector.”



MARS

Serpent

RC6

Twofish

Rijndael

Advanced Encryption Standard Contest Finalists



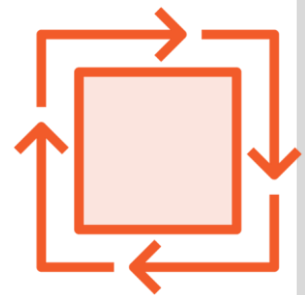
MARS



Block size – 128 bits plain text



Key size – 128 to 448 bits



Rounds - 32



Differentiation - MARS is not well suited for restricted-space environments due to its ROM requirement, which tends to be the highest among the finalists



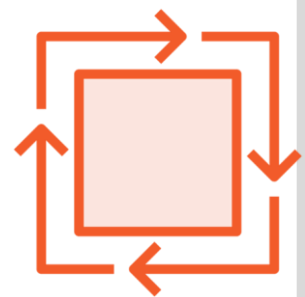
Serpent



Block size – 128 bits



Key size – 256 bits



Rounds - 32



Differentiation – of finalist slowest in software fastest in hardware processing



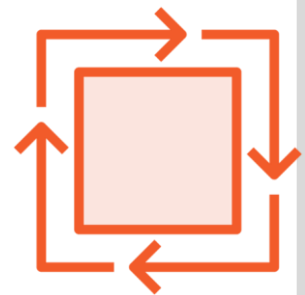
RC6



Block size – 128 bits



Key size – 256 bits



Rounds - 20



Differentiation - block, key, and round sizes are parameterized, it therefore supports key sizes much higher than 256 bits.



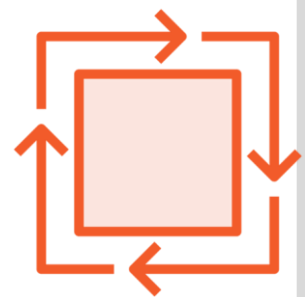
Twofish



Block size – 128 bits



Key size – 128, 192, 256 bits



Rounds - 16



Differentiation - throughput is somewhat reduced for the larger key sizes.



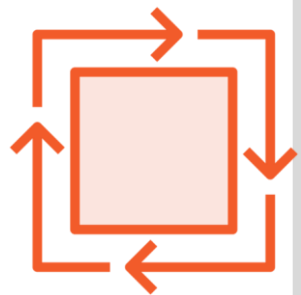
Rijndael



Block size – 128, 192, 256 bits



Key size - 128, 192, 256 bits



Rounds – 10, 12, 14



Differentiation -the key setup performance for Rijndael is consistently the fastest of all the finalists



“NIST selected Rijndael as the proposed AES algorithm at the end of a very long and complex evaluation process. During the evaluation, NIST analyzed all public comments, papers, verbal comments at conferences, and NIST studies and reports. NIST judged Rijndael to be the best overall algorithm for the AES.”

Report on the Development of the Advanced Encryption Standard (AES)–
May-June 2001 - NIST



Demo

Let's look at a symmetric key and cryptosystem's actions on plaintext

- This will help to understand the application of cryptography on plaintext
- We will use CrypTool to demonstrate confidentiality and access control



Asymmetric Algorithms Overview



Five Rules of Asymmetric Encryption

Process Order

When one half of key-pair encrypts the other decrypts

Public Key

Encryption objective is confidentiality and access control

Private Key

Encryption objective is integrity, authenticity, and non-repudiation

Digital Signature

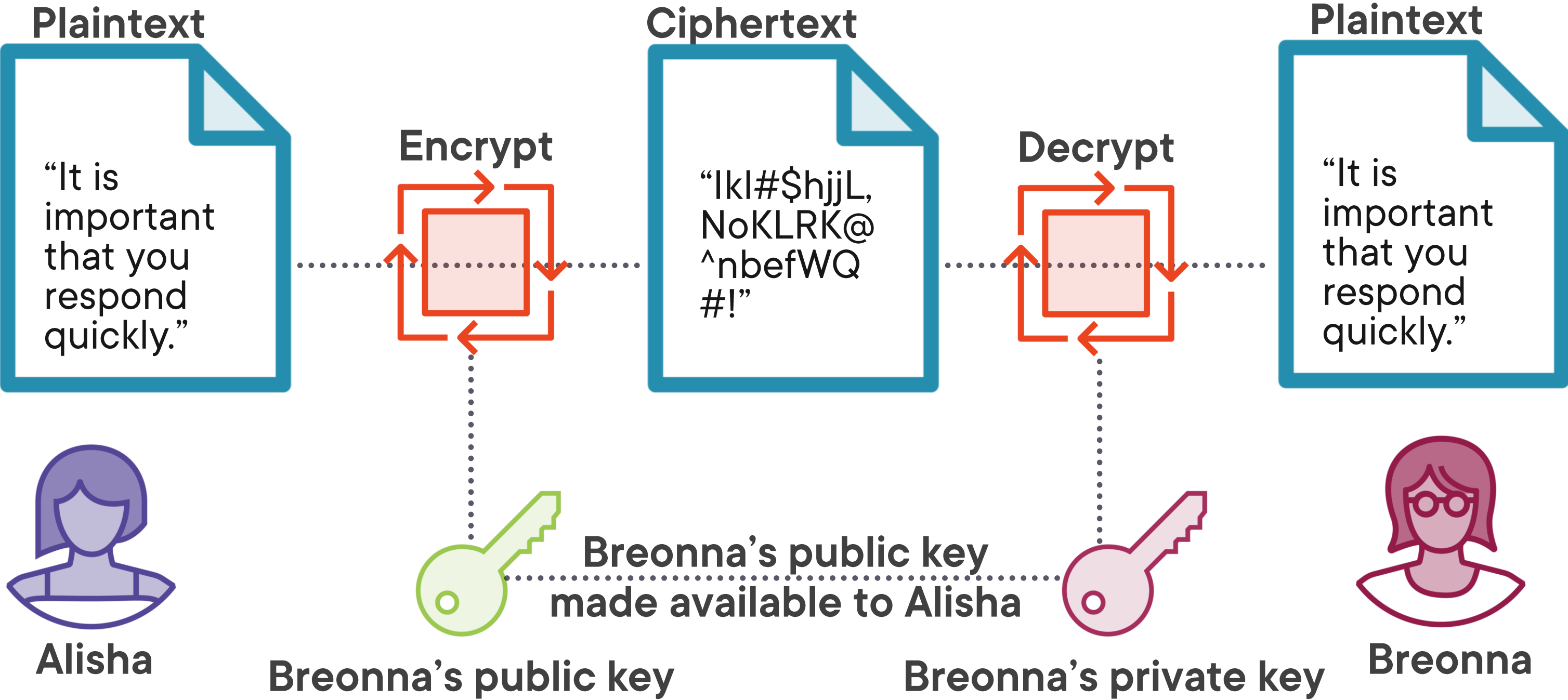
Private key (encrypting) signing a digest

Digital Certificate

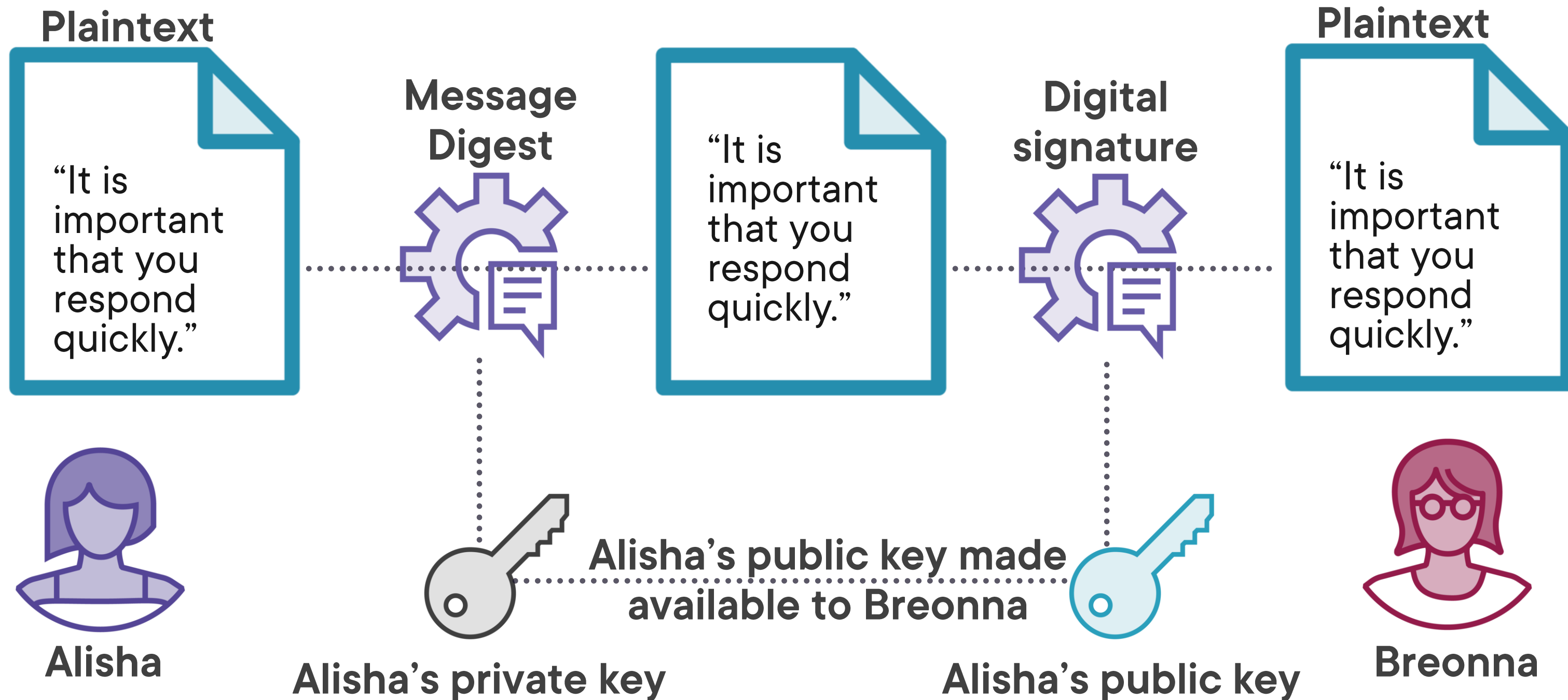
Digital document containing DS of CA and public key of owner



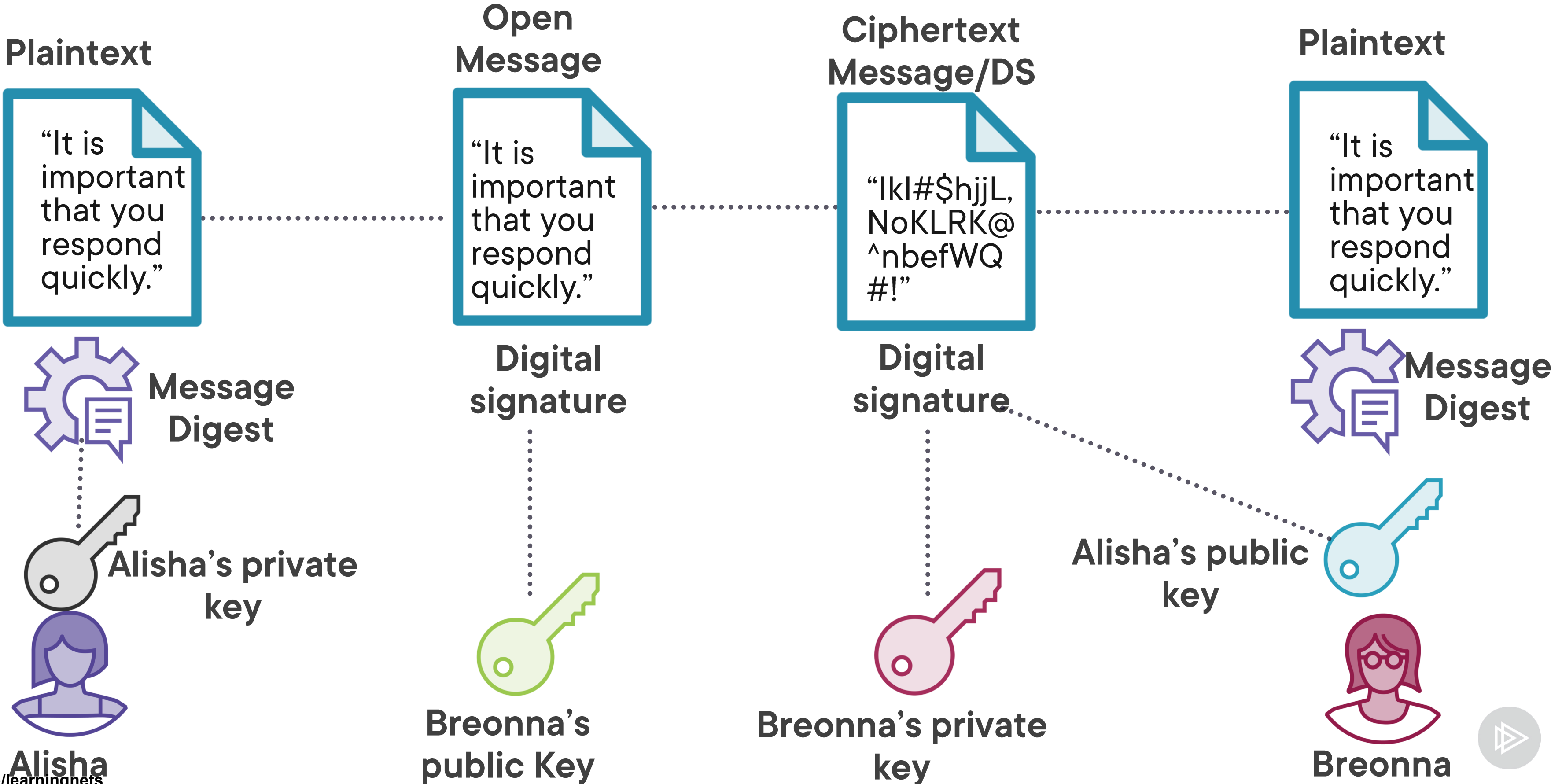
Asymmetric Key Process Flow for Confidentiality



Asymmetric Key Process Flow for Non-repudiation



Asymmetric Key Process Flow for Non-repudiation and Confidentiality



The Initial Asymmetric Algorithm



“Public key distribution systems offer a different approach to eliminating the need for a secure key distribution channel. In such a system, two users who wish to exchange a key communicate back and forth until they arrive at a key in common. A third-party eavesdropping on this exchange must find it computationally infeasible to compute the key from the information overheard.”

New Directions in Cryptography – November 1976 – IEEE vol IT 22
Whitfield Diffie and Martin E. Hellman



Diffie-Hellman-Merkle Modulus Math Problem

Publicly Accessible Numbers

$$G=11 \quad P=17$$



Alisha



Breonna



Diffie-Hellman-Merkle Modulus Math Problem

Publicly Accessible Numbers

$$G=11 \ P=17$$

$$G=11 \ P=17$$



Alisha



Breonna



Diffie-Hellman-Merkle Modulus Math Problem

Publicly Accessible Numbers

$G=11$ $P=17$



Alisha

$G=11$ $P=17$



Breonna



Diffie-Hellman-Merkle Modulus Math Problem

Privately retained numbers
independently selected by each

5

$G=11$ $P=17$



Alisha

9

$G=11$ $P=17$



Breonna



Diffie-Hellman-Merkle Modulus Math Problem

$$11^5 \bmod 17 = 10$$

5 remains a secret

$$G=11 \quad P=17$$



Alisha

9

$$G=11 \quad P=17$$



Breonna

10



Diffie-Hellman-Merkle Modulus Math Problem

$$11^5 \text{ mod } 17 = 10$$

5 remains a secret

$$G=11 \quad P=17$$



Alisha

6

$$11^9 \text{ mod } 17 = 6$$

9 remains a secret

$$G=11 \quad P=17$$



Breonna

10



Diffie-Hellman-Merkle Modulus Math Problem

$$6^5 \bmod 17 = 7$$

$$11^5 \bmod 17 =$$

$$G=11 \quad P=17$$



Alisha

$$11^9 \bmod 17 =$$

9 remains a secret

$$G=11 \quad P=17$$



10

Breonna



Diffie-Hellman-Merkle Modulus Math Problem

$$6^5 \bmod 17 = 7$$

$$11^5 \bmod 17 =$$

$$G=11 \quad P=17$$



Alisha

$$10^9 \bmod 17 = 7$$

$$11^9 \bmod 17 =$$

$$G=11 \quad P=17$$



Breonna



Diffie-Hellman-Merkle Modulus Math Problem

$$6^5 \bmod 17 = 7$$



Shared-secret is 7



$$10^9 \bmod 17 = 7$$

$$11^5 \bmod 17 =$$

$$11^9 \bmod 17 =$$

G=11 P=17

G=11 P=17



Alisha



Breonna



Asymmetric Algorithm Types and Characteristics



Diffie-Hellman-Merkle



Primary function – negotiation “exchange” of symmetric keys



Primary mathematics – discrete logarithms over finite fields



Distinguishing characteristic – first commercially viable asymmetric algorithm



RSA



Primary function – session keys, digital signatures, and message confidentiality



Primary mathematics – factoring the product of two large prime numbers



Distinguishing characteristic – most widely used asymmetric algorithm in history



ElGamal



Primary function – session keys, digital signatures, and message confidentiality



Primary mathematics – discrete logarithms over finite fields



Distinguishing characteristic – used concepts of Diffie-Hellman-Merkle for key distribution while introducing digital signature scheme



Elliptic Curve Cryptography (ECC)



Primary function – session keys, digital signatures, and message confidentiality



Primary mathematics – algebraic structure of elliptic curves over finite fields



Distinguishing characteristic – shorter key lengths uses less computational power



ECC vs. RSA

ECC Key length (bits)

160

224

256

384

512

RSA Key length (bits)

1024

2048

3072

7680

15360



Asymmetric cryptography is too slow.



Demo

Let's generate a private/public key pair

- This will allow confidential protection of host device connecting remotely and non-repudiation
- We will use CLI feature of SSH from a command prompt



Hashing Algorithms Overview



Parity bits

Checksums

**Cyclic Redundancy
Checks (CRC)**

**Message authentication
codes (MAC and HMAC)**

Hashing

System Integrity Codes



**Easy to compute the hash
value for a message**

**Infeasible to generate a
message given hash**

**Infeasible to modify a
message without
changing hash**

**Difficult to find two
different messages with
the same hash**

Hashing Primitives



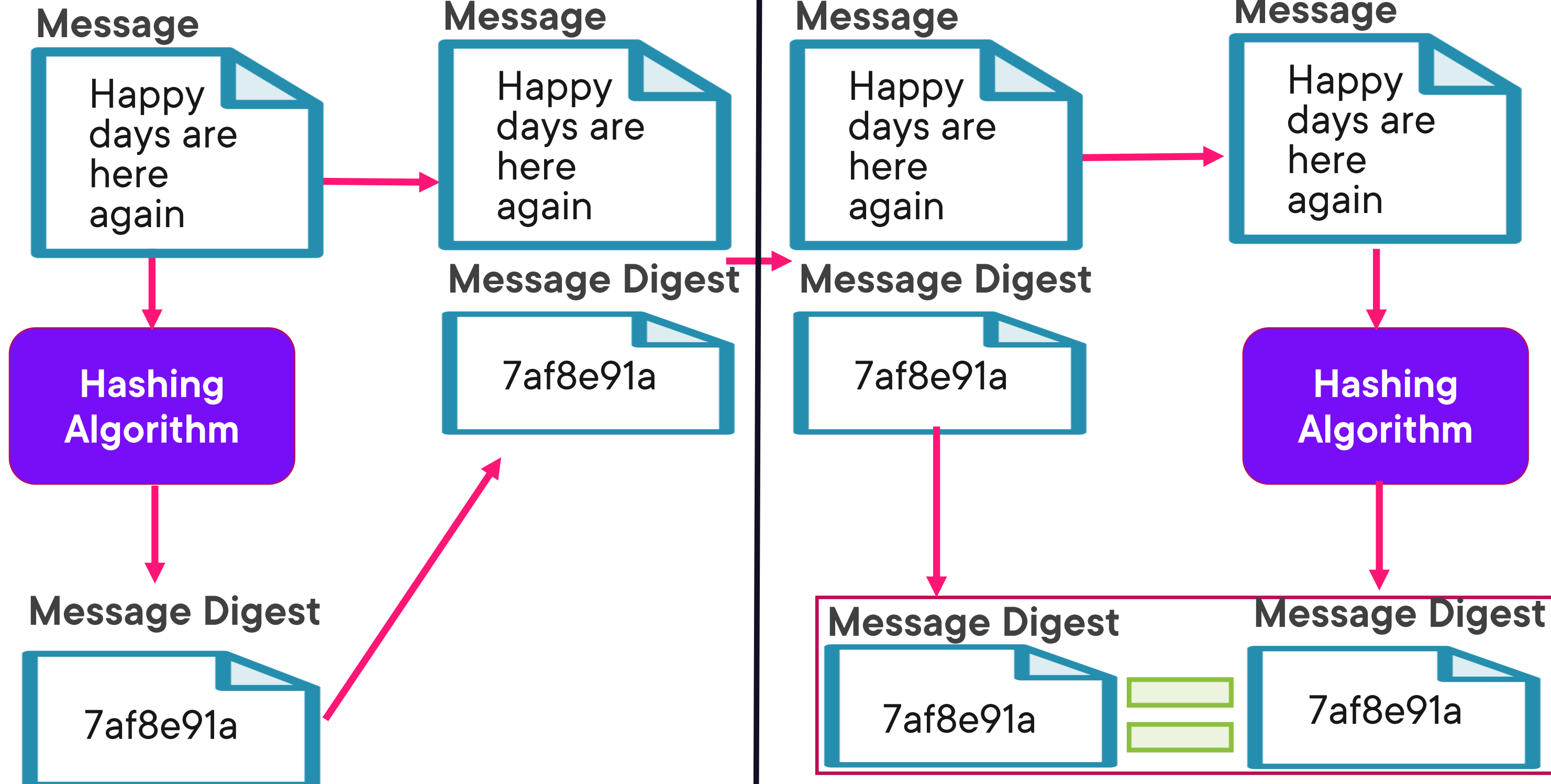


Alisha

Hash Operation



Breonna



Digest Sensitivity to Change

The digest of this text is below in green. It was produced with SHA256

```
a8cd0052dc0da24  
9082747c83bf5ac2  
9d7246f648b1981c2  
c91e0e7a0d0eba77
```

The digest of this text is below in green. It was produced with SHA256.

```
a686c6b46079323e  
41e7e18555b0a3d4  
d47b6d26016d11df2  
097c83915312e69
```

The digest of this text is below in green. It was produced with SHA256!

```
d623ae37719173936  
cc788dc325863774  
adb5d1822be8d4ab  
4cdd354971bb0b0
```



Hashing Algorithm Types and Characteristics



MD2

MD4

MD5

Message Digest (MD)



SHA-0

SHA-1

SHA-2

SHA-3

Secure Hash Algorithm (SHA)



Additional Hash Algorithms

HAVAL

**128-bit block, variable-bit
digest**

RIPEMD-160

512-bit block, 160-bit digest



Up Next:

Up Next:

Secure Protocols and Cryptographic Lifecycles

